

R-SETS AND CATEGORY

BY

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ABSTRACT. We prove some category theoretic results for R -sets much in the spirit of Vaught and Burgess. Since the proofs entail many results on R -sets and the R -operator, we have studied them in some detail and have formulated many results appropriate for our purpose in, perhaps, a more unified manner than is available in the literature. Our main theorem is the following: Any R -set in the product of two Polish spaces can be approximated, in category, uniformly over all sections, by sets generated by rectangles with one side an R -set and the other a Borel set. In fact, we prove a levelwise version of this result. For C -sets, this has been proved by V. V. Srivatsa.

1. Introduction and preliminaries. The theory of R -sets and the R -operator, introduced by Kolmogorov almost half a century ago, has been studied extensively by Russian mathematicians [10, 12, 13] and most of the basic properties have been deduced by them. However, it is only very recently that interest in the theory has been revived due mainly to the work of Hinman [7, 9] who developed the effective counterpart and showed that the effective hierarchies have deep interconnections with recursion-theoretic hierarchies. The introduction of Borel-programmable (BP) sets by Blackwell added a new dimension to the theory, and since then it has been shown by Burgess and Lockhart that the hierarchy obtained from BP-sets by iteration gives precisely the R -sets [6]. That two seemingly different definitions yield the same class of sets suggests that the R -sets form a natural class of subsets of the reals. Burgess has also obtained a different characterization for R -sets. He has proved that the entire hierarchy of R -sets can be obtained by applying the game quantifier to the "difference hierarchy" (of Δ_3^0) obtained from sequences of G_δ -sets [3]. Hinman (and also independently Aczel and Vaught, for the first level) first observed that the theory of inductive definability and games can be effectively used to study the hierarchy of R -sets [8, 1, 19]. These are the major tools employed by Burgess to prove most of his results on R -sets [3, 4]. In this paper, to obtain our main result, viz. the approximation theorem, we have taken recourse to these methods.

R -sets in $X \times Y$, the product of two Polish spaces, are in general complicated sets and cannot be related to any reasonable product σ -field. For instance, as observed by B. V. Rao [16], C -sets in $\mathbb{R} \times \mathbb{R}$ need not belong to $\mathcal{L}(\mathbb{R}) \otimes \mathcal{L}(\mathbb{R})$, the product of the σ -fields of Lebesgue measurable sets. However, as shown by V. V. Srivatsa [17],

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C -sets in $X \times Y$ can be “approximated” sectionwise, in the sense of measure and category, by sets in product σ -fields. His methods do not seem to generalize to R -sets. In this paper, we have shown that R -sets can be similarly approximated, in category, by sets in product σ -fields. This incidentally gives the selection theorem of Burgess (cf. [3]).

To obtain the approximation theorem, it is necessary to reprove some of Burgess’ results in a more rigorous, accessible and, perhaps, transparent style. Our paper, therefore, serves a two-fold purpose.

The paper is organized as follows. Positive analytical operations and $\delta - s$ operations are introduced in §2, and some of their elementary properties are discussed there. The papers of Kantorovitch and Livenson [10] give a detailed exposition of these operations. In §3, we have studied the operator R of Kolmogorov and the more general operator \mathcal{R}_N . In this section, we have also shown how inductive definability and games arise in the context of these operations. The theory of R -sets, studied in great detail by Lyapunov [12, 13], is treated in §4. In this section, we have given a proof of the pre-well-ordering property enjoyed by the classes $c\mathcal{R}^\rho$ ($\rho < \omega_1$) via the comparison of indices lemma. The proof of this lemma is much along the lines of the Kunen-Aczel theorem (cf. [14]). The comparison of indices lemma is crucial for our purpose since it helps in computing the complexity of the winning strategy for the existential player in the game associated with the operator R . This is done in §5. Here we have also obtained a decomposition of $E^* = \{x : E^x \text{ is comeager}\}$ for sets $E \in \Sigma_1^{\Phi^*}$, analogous to the one obtained by Vaught for analytic sets [19] and for \mathcal{R}^1 -sets obtained by Burgess and Miller [5]. This immediately gives us the transfer theorem (cf. 5.4), which essentially computes E^* when E is computed by $R\Phi_N$, whenever the computation for F^* for sets F computed by Φ_N is known. §6 deals with a few applications of the transfer theorem, viz. the computation of E^* for R -sets E . It is worth mentioning here that although our methods for computing E^* are implicit in Burgess’ proof for the same (cf. [5]), he computes E^* only for \mathcal{R}^1 -sets E , as any computation for higher levels using his methods will involve great notational complexities. By restricting ourselves to certain games of length ω and isolating the “core” of his proof (viz. the transfer theorem) we have been able to compute E^* for all levels of the hierarchy of R -sets by a simple inductive argument. These computations yield sets in product σ -fields which “approximate” sectionwise (in the sense of category) R -sets in two dimensions. This is done in §7. Incidentally, in this section, we have proved a slightly stronger version of the Game Formula of Kechris (cf. [11]) needed for our purpose.

For our notation and terminology we shall mainly follow Moschovakis [15]. The letter ω will denote the set of natural numbers and ω^ω the set of all sequences of natural numbers equipped with the product of discrete topologies. Letters $\alpha, \beta, \gamma, \delta, \dots$ will serve as variables over ω^ω and η, ξ, \dots as variables over 2^ω . Seq will denote the set of sequence numbers of all finite sequences of natural numbers. We will mainly use s, t, u, v to denote sequence numbers. We fix a base $\Sigma(s)$ for the topology of ω^ω , where

$$\Sigma(s) = \{ \alpha \in \omega^\omega : \bar{\alpha}(\text{lh}(s)) = s \}.$$

If s and t are sequence numbers, we write $s < t$ if $s = t \upharpoonright i$ for some $i < \text{lh}(s)$; $s * t$ or $\hat{s}t$ is the catenation of s and t . If $s = \langle a_0, \dots, a_{k-1} \rangle$, then $(s)_i = a_i$ for $i < k$. e or $\langle \rangle$ will denote the empty sequence as well as its code.

If \mathcal{F} is a collection of subsets of X , then $\sigma(\mathcal{F})$ denotes the σ -field generated by \mathcal{F} . If \mathcal{B} and \mathcal{C} are σ -fields on T and X , respectively, $\mathcal{B} \otimes \mathcal{C}$ denotes the product σ -field. For X a separable metric, \mathcal{B}_X denotes its Borel σ -field.

Given a monotone set relation $\Gamma(w, x, A)$, where w varies over W and A varies over subsets of W , Γ_x denotes the induced set operation

$$\Gamma_x(A) = \{w \in W : \Gamma(w, x, A)\}; \quad x \in X.$$

Γ_x^μ denotes the μ th iterate, viz.

$$\Gamma_x^\mu = \Gamma_x \left(\bigcup_{\nu < \mu} \Gamma_x^\nu \right).$$

We define

$$\Gamma_x^\infty = \bigcup_{\mu} \Gamma_x^\mu$$

and put

$$w \in \Gamma_x^{<\mu} \leftrightarrow (\exists \nu < \mu) [w \in \Gamma_x^\nu].$$

The fixed point of Γ is

$$\Gamma^\infty = \{ (w, x) : w \in \Gamma_x^\infty \}$$

and is called the relation built up by Γ . We shall use elementary facts on inductive definability and games as found in [15]. Note that all unexplained notation and terminology is from Moschovakis [15].

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2. Positive analytical operations. In this section, we shall discuss positive analytical operations and some of their properties needed for our purpose.

Let X be a nonempty set and $N \subseteq \omega^\omega$. Let $\{E_n : n \in \omega\}$ be a sequence of subsets of X .

DEFINITION 2.1. The $\delta - s$ operation with base N is defined by

$$\Phi_N(\{E_n : n \in \omega\}) = \bigcup_{\alpha \in N} \bigcap_{k=0}^{\infty} E_{\alpha(k)}.$$

In most cases we shall take the base set $N \subseteq 2^\omega$ so that

$$\Phi_N(\{E_n : n \in \omega\}) = \bigcup_{\eta \in N} \bigcap_{n \in \eta} E_n.$$

To avoid trivialities, we shall always assume $\emptyset \notin N$ and $N \neq \emptyset$.

EXAMPLES. If $N = \{\lambda n \bar{\alpha}(n) : \alpha \in \omega^\omega\}$, then $\Phi_N = \text{operation } \mathcal{A}$. If

$$N = \{\alpha \in \omega^\omega \mid \text{Range}(\alpha) \text{ is infinite}\},$$

then $\Phi_N = \text{lim sup}$. If $N = \{(n, n, \dots) \mid n \in \omega\}$, then $\Phi_N = \bigcup$ (countable).

An operation over X is a function $\Phi : \mathcal{P}(X)^\omega \rightarrow \mathcal{P}(X)$.

DEFINITION 2.2. An operation over X is said to be a positive analytical operation if

(a) Φ is nonconstant and

(b) $x \in \Phi(\langle E_n : n \in \omega \rangle) \& y \notin \Phi(\langle F_n : n \in \omega \rangle) \rightarrow (\exists n)(x \in E_n \& y \notin F_n)$.

Clearly, \cup , \mathcal{A} , \limsup are positive analytical operations. A positive analytical operation on constant sequences takes on the constant value and is isotone i.e. for any families $\{F_n : n \in \omega\}$ and $\{G_n : n \in \omega\}$, if $(\forall n)(F_n \subseteq G_n)$, then $\Phi(\langle F_n : n \in \omega \rangle) \subseteq \Phi(\langle G_n : n \in \omega \rangle)$.

Given a positive analytical operation Φ over X , one can define a positive analytical operation Φ' over any set Y as follows. For any family $\{F_n : n \in \omega\}$ of subsets of Y and $y \in Y$, put

$$y \in \Phi'(\langle F_n : n \in \omega \rangle) \leftrightarrow \text{there is a family } \{E_n : n \in \omega\} \text{ of subsets of } X \\ \text{and } x \in X \text{ such that } x \in \Phi(\langle E_n : n \in \omega \rangle) \\ \text{and } (\forall n)(x \in E_n \rightarrow y \in F_n).$$

It is easy to check that Φ' is positive analytical and $\Phi' = \Phi$ when $Y = X$. The operation Φ' is called the extension of Φ over Y . Henceforth we shall use the same symbol Φ to denote a positive analytical operation together with all its extensions.

Notice that a $\delta - s$ operation is a positive analytical operation (over any set X). The converse is true and follows from the following.

PROPOSITION. Let Φ be a positive analytical operation over X . Then $\Phi = \Phi_N$ for some $N \subseteq 2^\omega$ ($N \subseteq \omega^\omega$).

PROOF. Let

$$D_i = \{ \eta \in 2^\omega \mid i \in \eta \}, \quad \text{or,} \\ = \{ \alpha \in \omega^\omega \mid i \in \text{Range}(\alpha) \}, \quad i \in \omega.$$

Put $N = \Phi(\langle D_i : i \in \omega \rangle)$. It is easy to check that $\Phi = \Phi_N$. The set N obtained above is called the canonical base for Φ .

DEFINITION 2.3. A base $N \subseteq \omega^\omega$ is complete if $N = \Phi_N(\langle D_n : n \in \omega \rangle)$, when D_n is defined as in the Proposition. Thus the canonical base for a positive analytical operation is complete.

Equivalently, a base $N \subseteq 2^\omega$ is complete if

$$\eta \in N \& \eta \subseteq \eta' \subseteq \omega \rightarrow \eta' \in N.$$

If $\tilde{N} = \Phi_N(\langle D_n : n \in \omega \rangle)$, then \tilde{N} is complete (called the completion of N) and, moreover, $\Phi_{\tilde{N}} = \Phi_N$.

DEFINITION 2.4. For any operation Φ , the dual Φ^0 is defined by

$$\Phi^0(\langle E_n : n \in \omega \rangle) = [\Phi_N(\langle E_n^c : n \in \omega \rangle)]^c$$

e.g. $\cup^0 = \cap$, $(\limsup)^0 = \liminf$, $\mathcal{A}^0 = \Gamma$, where

$$\Gamma(\langle E_n : n \in \omega \rangle) = \{ x : (\forall \alpha)(\exists n)(x \in E_{\bar{\alpha}(n)}) \}.$$

If Φ_N is a $\delta - s$ operation with base $N \subseteq 2^\omega$, then the canonical base N^0 of its dual is given by

$$(1) \quad \begin{aligned} N^0 &= \{ \eta \in 2^\omega : \eta \cap \eta' \neq \emptyset \text{ for every } \eta' \in N \} \\ &= \{ \eta \in 2^\omega : \eta^c \notin N \}, \text{ if } N \text{ is complete.} \end{aligned}$$

Plainly, N^0 is always complete. Thus, for any family $\{ E_n : n \in \omega \}$,

$$(2) \quad (\exists \eta \in N^0)(\forall n \in \eta)[x \in E_n] \leftrightarrow (\forall \eta \in N)(\exists n \in \eta)[x \in E_n].$$

If N is complete, $N^{00} = N$ and hence

$$(3) \quad (\forall \eta \in N^0)(\exists n \in \eta)[x \in E_n] \leftrightarrow (\exists \eta \in N)(\forall n \in \eta)[x \in E_n].$$

3. The R -operator. Although the R -operator was first introduced by Kolmogorov, the first published account of the theory appeared in [10] and further results obtained in [12, 13]. Lyapunov also studied the hierarchies of R -sets (cf. §4) and obtained most of their properties. The interconnection between R -operators and games was first noticed by Hinman [8] (and also independently by Aczel [1]). Hinman also developed the effective theory and did most of the groundwork. Much of the material in this section is adapted from these sources.

DEFINITION 3.1. Let $\mathcal{N} = \{ N_p \subseteq 2^\omega : p \in \omega \}$ be a sequence of nonempty bases. $\theta \subseteq \omega$ is called an \mathcal{N} -chain if

- (a) $e \in \theta$,
- (b) $s \in \theta$ and $t < s \rightarrow t \in \theta$,
- (c) $s \in \theta \rightarrow \{ n : \hat{s}\langle n \rangle \in \theta \} \in N_s$.

Put $\Theta_{\mathcal{N}} = \{ \theta : \theta \text{ is an } \mathcal{N}\text{-chain} \}$.

$\mathcal{R}_{\mathcal{N}}$ is the set operation defined by

$$\mathcal{R}_{\mathcal{N}}(\{ E_n : n \in \omega \}) = \bigcup_{\theta \in \Theta_{\mathcal{N}}} \bigcap_{s \in \theta} E_s.$$

Clearly, $\mathcal{R}_{\mathcal{N}}$ is a $\delta - s$ operation with base $\Theta_{\mathcal{N}}$.

If $N_p = N$ for each p , then $\mathcal{R}_{\mathcal{N}}$ is denoted by $R\Phi_N$ and its base by RN . An \mathcal{N} -chain will then be called an N -chain.

EXAMPLES. Let $N = \{ \langle n \rangle : n \in \omega \}$ so that $\Phi_N = \cup$ and put $\mathcal{N} = \{ N \}$. Clearly, an \mathcal{N} -chain is any set of the form $\{ \bar{\alpha}(n) : n \in \omega \}$. Thus $R\cup = R\Phi_N = \mathcal{A}$. If $N = \{ \omega \}$, then $\Phi_N = \cap$ (countable) and the only \mathcal{N} -chain is ω , so $R\cap = \cap$.

DEFINITION 3.2. Let Φ_N and Φ_M be two $\delta - s$ operations with bases $N, M \subseteq 2^\omega$. The composed operation Ψ is given by

$$\Psi(\{ F_n : n \in \omega \}) = \Phi_N(\{ \Phi_M(\{ F_{\langle p, n \rangle} : n \in \omega \}) : p \in \omega \}).$$

Ψ is sometimes denoted by $\Phi_N\Phi_M$. By the characterization lemma, Ψ is a $\delta - s$ operation whose canonical base we shall denote by NM . Thus $\Psi = \Phi_N\Phi_M = \Phi_{NM}$ and

$$\eta \in NM \leftrightarrow (\exists \eta_1 \in N)(\forall n_1 \in \eta_1)(\exists \xi_1 \in M)(\forall m_1 \in \xi_1)[\langle n_1, m_1 \rangle \in \eta].$$

Henceforth, for simplicity, we shall take $X = \omega^\omega$ or $(\omega^\omega)^k \times \omega^1$, although most of the results hold for a general Polish space.

DEFINITION 3.3. For any operation Φ , let Σ_1^Φ be the class of relations of the form $\Phi(\langle F_n : n \in \omega \rangle)$ with all F_n clopen, Π_1^Φ the class of complements of such relations, and $\Delta_1^\Phi = \Sigma_1^\Phi \cap \Pi_1^\Phi$. Then $\Sigma_1^\cup = \Sigma_1^0$ and $\Sigma_1^\cap = \Sigma_1^1$.

If $\Phi = \Phi_N$, then define $\Phi^* = R\Phi_{NN^0}$.

The next two lemmas show a close connection between R -operators, inductive definability and games.

LEMMA 3.4. (a) Suppose $F = R\Phi_N(\langle F_n : n \in \omega \rangle)$. Then

$$(*) \quad x \in F \leftrightarrow (\exists \eta_0 \in N)(\forall n_0 \in \eta_0)(\exists \eta_1 \in N)(\forall n_1 \in \eta_1) \cdots \\ \cdots (\forall k) [x \in F_{\langle n_0, n_1, \dots, n_{k-1} \rangle}].$$

(b) If $E = \Phi_N^*(\langle E_n : n \in \omega \rangle)$, then

$$x \in E \leftrightarrow (\exists \eta_0 \in N)(\forall n_0 \in \eta_0)(\forall \xi_0 \in N)(\exists m_0 \in \xi_0)(\exists \eta_1 \in N)(\forall n_1 \in \eta_1) \\ (\forall \xi_1 \in N)(\exists m_1 \in \xi_1) \cdots (\forall k) [x \in E_{\langle \langle n_0, m_0 \rangle, \dots, \langle n_{k-1}, m_{k-1} \rangle \rangle}].$$

(The right-hand side of each equivalence is interpreted in terms of games between two players \forall and \exists .)

PROOF. Clearly (b) follows from (a) and the fact that

$$\eta \in NN^0 \leftrightarrow (\exists \eta_1 \in N)(\forall n_1 \in \eta_1)(\forall \xi_1 \in N)(\exists m_1 \in \xi_1) [\langle n_1, m_1 \rangle \in \eta].$$

To prove the first assertion, fix x and suppose $x \in F$. Get an N -chain $\theta \in RN$ such that $(\forall s \in \theta)[x \in F_s]$. Now, consider the following strategy for \exists . As his first move \exists plays $\eta_0 = \{n : \langle n \rangle \in \theta\}$ which is clearly in N . Any response $n_0 \in \eta_0$ by \forall gives a set $\eta_1 = \{n : \langle n_0, n \rangle \in \theta\} \in N$ and \exists should play η_1 as his next move. If \forall then plays $n_1 \in \eta_1$, we still have $\langle n_0, n_1 \rangle \in \theta$ and \exists responds with $\eta_2 = \{n : \langle n_0, n_1, n \rangle \in \theta\}$. If \exists follows this strategy, then clearly for any k , $\langle n_0, n_1, \dots, n_{k-1} \rangle \in \theta$ and so $x \in F_{\langle n_0, n_1, \dots, n_{k-1} \rangle}$. Hence it is a winning strategy for \exists in the game $(*)$.

For the converse implication suppose σ is a winning strategy for \exists . Let θ be the set of sequences $\langle n_0, \dots, n_{k-1} \rangle$ of first k possible moves of player \forall when \exists follows this strategy σ . Clearly θ is an N -chain and since σ is a winning strategy for \exists ,

$$(\forall k) [\langle n_0, \dots, n_{k-1} \rangle \in \theta \rightarrow x \in F_{\langle n_0, n_1, \dots, n_{k-1} \rangle}].$$

Hence $x \in R\Phi_N(\langle F_n : n \in \omega \rangle) = F$.

REMARK. It follows from Lemma 3.4 that our definition of Φ^* is equivalent (cf. Definition 3.9) to that introduced in [8, V. 4].

The next result is due to Hinman [8].

THEOREM 3.5. For any positive analytical operation Φ and any $E \subseteq (\omega^\omega)^k \times \omega^1$ in $\Pi_1^{\Phi^*}$, there exists a (monotone) inductive operator Γ such that for all x

$$E(x) \leftrightarrow e \in \Gamma_x^\infty.$$

PROOF. Let $\{E_s : s \in \omega\}$ be a family of clopen subsets of $(\omega^\omega)^k \times \omega^1$ such that

$$x \notin E \leftrightarrow x \in \Phi^*(\{E_s : s \in \omega\}).$$

Let $N \subseteq 2^\omega$ be the canonical base for Φ . Then $\Phi^* = R\Phi_{NN^0}$. Define a set relation operative on ω as follows:

(4)

$$s \in \Gamma_x(A) \leftrightarrow x \notin E_s \vee (\forall \eta \in N)(\exists n \in \eta)(\exists \xi \in N)(\forall m \in \xi)[s * \langle \langle n, m \rangle \rangle \in A].$$

Clearly, $\Gamma(s, x, A)$ is a monotone set relation. Put $E^s = \Phi^*(\langle E_{s \cdot t} : t \in \omega \rangle)$. One can easily see that $E^e = E^c$ and $E^s \subseteq E_s$. We claim that for all s ,

(5)

$$x \notin E^s \leftrightarrow s \in \Gamma_x^\infty$$

and the result follows by putting $s = \langle \rangle$. We shall prove the implication (\leftarrow) by induction. If $s \in \Gamma_x^0$, then $x \notin E_s \supseteq E^s$. Now suppose $s \in \Gamma_x^\mu$, $\mu > 0$. Then

$$x \notin E_s \vee (\forall \eta \in N)(\exists n \in \eta)(\exists \xi \in N)(\forall m \in \xi)[s * \langle \langle n, m \rangle \rangle \in \Gamma_x^{<\mu}].$$

If $x \notin E_s$, we are done; otherwise by induction hypothesis

$$(\forall \eta \in N)(\exists n \in \eta)(\exists \xi \in N)(\forall m \in \xi)[x \notin E^{s * \langle \langle n, m \rangle \rangle}]$$

which by Lemma 3.4 (and determinacy) implies

$$\begin{aligned} & (\forall \eta \in N)(\exists n \in \eta)(\exists \xi \in N)(\forall m \in \xi) \{ (\forall \eta_1 \in N)(\exists n_1 \in \eta_1) \\ & (\exists \xi_1 \in N)(\forall m_1 \in \xi_1) \cdots (\exists k) [x \notin E_{s * \langle \langle n, m \rangle \rangle * \langle \langle n_1, m_1 \rangle \rangle \cdots \langle \langle n_k, m_k \rangle \rangle}] \}. \end{aligned}$$

This clearly implies

$$\begin{aligned} & (\forall \eta_0 \in N)(\exists n_0 \in \eta_0)(\exists \xi_0 \in N)(\forall m_0 \in \xi_0)(\forall \eta_1 \in N)(\exists n_1 \in \eta_1) \\ & (\exists \xi_1 \in N)(\forall m_1 \in \xi_1) \cdots (\exists k) [x \notin E_{s * \langle \langle n_0, m_0 \rangle \rangle \cdots \langle \langle n_{k-1}, m_{k-1} \rangle \rangle}] \end{aligned}$$

and thus $x \notin R\Phi_{NN^0}(\langle E_{s \cdot t} : t \in \omega \rangle) = E^s$, by Lemma 3.4 again.

Conversely, let $s \notin \Gamma_x^\infty$. We shall show that $x \in E^s$ i.e.,

$$\begin{aligned} (*) \quad & (\exists \eta_0 \in N)(\forall n_0 \in \eta_0)(\forall \xi_0 \in N)(\exists m_0 \in \xi_0)(\exists \eta_1 \in N)(\forall n_1 \in \eta_1) \\ & (\forall \xi_1 \in N)(\exists m_1 \in \xi_1) \cdots (\forall k) [x \in E_{s * \langle \langle n_0, m_0 \rangle \rangle \cdots \langle \langle n_{k-1}, m_{k-1} \rangle \rangle}]. \end{aligned}$$

Since $s \notin \Gamma_x(\Gamma_x^\infty)$, by definition of Γ , $x \in E_s$ and moreover,

$$(\exists \eta_0 \in N)(\forall n_0 \in \eta_0)(\forall \xi_0 \in N)(\exists m_0 \in \xi_0)[s * \langle \langle n_0, m_0 \rangle \rangle \notin \Gamma_x^\infty].$$

Now, \exists can win the game $(*)$ by adopting the following strategy. He picks $\eta_0 \in N$ such that for any choice of $n_0 \in \eta_0$ and $\xi_0 \in N$, there is an $m_0 \in \xi_0$ such that $s * \langle \langle n_0, m_0 \rangle \rangle \notin \Gamma_x^\infty = \Gamma_x(\Gamma_x^\infty)$. Thus $x \in E_{s * \langle \langle n_0, m_0 \rangle \rangle}$ and

$$(\exists \eta_1 \in N)(\forall n_1 \in \eta_1)(\forall \xi_1 \in N)(\exists m_1 \in \xi_1)[s * \langle \langle n_0, m_0 \rangle \rangle, \langle \langle n_1, m_1 \rangle \rangle \notin \Gamma_x^\infty].$$

\exists then picks $\eta_1 \in N$ such that for any choice of $n_1 \in \eta_1$ and $\xi_1 \in N$ made by \forall , there is an $m_1 \in \xi_1$ (and \exists plays such an m_1) such that $s * \langle \langle n_0, m_0 \rangle \rangle, \langle \langle n_1, m_1 \rangle \rangle \notin \Gamma_x^\infty$. Proceeding this way, \exists has a strategy which ensures that for all k , $x \in E_{s * \langle \langle n_0, m_0 \rangle \rangle \cdots \langle \langle n_{k-1}, m_{k-1} \rangle \rangle}$ and so \exists wins the game $(*)$. Consequently, $x \in E^s$.

REMARK. If E is such that $x \notin E \leftrightarrow x \in R\Phi_N(\langle E_s : s \in \omega \rangle)$, then the inductive operator takes a simpler form, viz.,

$$(6) \quad s \in \Gamma_x(A) \leftrightarrow x \notin E_s \vee (\forall \eta \in N)(\exists n \in \eta)[s * \langle n \rangle \in A].$$

More generally, if $\mathcal{N} = \{ N_p : p \in \omega \}$ is a sequence of bases and

$$E^c = \mathcal{R}_{\mathcal{N}}(\{ E_s : s \in \omega \}),$$

then we take the following inductive operator

$$(7) \quad s \in \Gamma_x(A) \leftrightarrow x \notin E_s \vee (\forall \eta \in N_s)(\exists n \in \eta)[s * \langle n \rangle \in A],$$

and the conclusion of the above theorem still holds.

The inductive operator (4) (or (6) or (7)) is called the canonical inductive operator associated with $\{ E_s : s \in \omega \}$ and N (or \mathcal{N}).

The above characterization theorem yields a decomposition of sets in $\Pi_1^{\Phi^*}(\Sigma_1^{\Phi^*})$ into simpler sets as is evident from the next theorem.

DEFINITION 3.6. For any operation Φ , $\nabla(\Phi)$ is the smallest class of relations containing clopen relations and closed under Φ and Φ^0 .

Thus $\nabla(\cup) = \nabla(\cap) = \Delta_1^1$.

Let $E \in \Sigma_1^{\Phi^*}$ and suppose $E = \Phi^*(\{ E_s : s \in \omega \})$. Let N and Γ be as above. Then by 3.5, $E(x) \leftrightarrow e \notin \Gamma_x^\infty$. Set

$$E_s^\mu = \{ x : s \notin \Gamma_x^\mu \}.$$

Then $E = \bigcap_{\mu < \omega_1} E_e^\mu$. Now define

$$T^\mu = \bigcup_{s \in \omega} (E_s^\mu - E_s^{\mu+1}).$$

It is easy to prove by induction on μ that for all $\mu < \omega_1$, $s \in \omega$, E_s^μ and T^μ are in $\nabla(\Phi)$. Then we have

$$\text{THEOREM 3.7. } E = \bigcup_{\mu < \omega_1} (E_e^\mu - T^\mu) = \bigcap_{\mu < \omega_1} E_e^\mu.$$

PROOF. Let $x \in E$. Define

$$\beta(s) = \begin{cases} \text{least ordinal } \rho \text{ such that } x \notin E_s^\rho & \text{if } (\exists \rho)(x \notin E_s^\rho), \\ 0 & \text{otherwise.} \end{cases}$$

Let $\omega_1 > \rho_0 > \beta(s)$, $\forall s$. Then $(\forall s)[x \in E_s^{\rho_0} \leftrightarrow x \in E_s^{\rho_0+1}]$. Consequently, $x \notin T^{\rho_0}$ and thus $x \in E_e^{\rho_0} - T^{\rho_0}$.

Conversely, suppose for some $\rho_0 < \omega_1$, $x \in E_e^{\rho_0} - T^{\rho_0}$. Then

$$(\forall s)[x \in E_s^{\rho_0} \leftrightarrow x \in E_s^{\rho_0+1}].$$

One can check by induction that

$$(\forall \rho > \rho_0)(\forall s)[x \in E_s^{\rho_0} \leftrightarrow x \in E_s^\rho].$$

So in particular, $x \in \bigcap_{\rho > \rho_0} E_e^\rho = \bigcap_{\rho < \omega_1} E_e^\rho = E$.

Standard arguments using the above decomposition and the countable chain condition yield the following (cf. [8]).

THEOREM 3.8. If Φ preserves measurability (preserves the Baire property), then so does Φ^* .

DEFINITION 3.9. For any two operations Φ and Ψ , Φ subsumes Ψ ($\Phi \geq \Psi$) if there is a function $f: \omega \rightarrow \omega$ such that for any family $\{ F_n : n \in \omega \}$,

$$\Psi(\{ F_n : n \in \omega \}) = \Phi(\{ F_{f(n)} : n \in \omega \}).$$

Φ and Ψ are said to be equivalent ($\Phi \sim \Psi$) if $\Phi \geq \Psi$ and $\Psi \geq \Phi$. For example, \mathcal{A} subsumes both \cup and \cap .

DEFINITION 3.10. A positive analytical operation Φ is said to be normal if there is a function g such that for any family $\{F_n : n \in \omega\}$,

$$\Phi(\{ \Phi(\{ F_{\langle p, q \rangle} : q \in \omega \}) : p \in \omega \}) = \Phi(\{ F_{g(n)} : n \in \omega \}).$$

We shall omit the proof of the next lemma which can be found in [7, 9].

LEMMA 3.11. For any operations Φ and Ψ

- (a) $\Phi^{00} = \Phi$;
- (b) $\Phi \geq \Psi \rightarrow \Phi^0 \geq \Psi^0$;
- (c) $\Phi \circ \Phi^0 \geq \Phi, \Phi^0$;
- (d) $R\Phi \geq \Phi$;
- (e) $\Phi \geq \Psi \rightarrow R\Phi \geq R\Psi$;
- (f) $R\Phi \geq \Psi$ and $R\Phi \geq \Psi^0 \rightarrow R\Phi \geq \Psi \cdot \Psi^0$;
- (g) $R\Phi \sim RR\Phi$;
- (h) $R\Phi$ is normal.

4. The R -sets. We shall first construct a sequence $\{R_\rho : \rho < \omega_1\}$ of positive analytical operations by induction as follows. Put $R_0 = \mathcal{A}$ and having defined R_ρ , put

$$R_{\rho+1} = R_\rho^*.$$

If λ is the limit, choose a sequence $\rho_i \uparrow \lambda$ and set, for any family $\{E_n : n \in \omega\}$,

$$\Phi(\{E_n : n \in \omega\}) = \bigcap_{i=0}^{\infty} \Phi_{N_{\rho_i} N_{\rho_i}^0}(\{E_{\langle i, m \rangle} : m \in \omega\}),$$

where N_{ρ_i} is the canonical base for R_{ρ_i} . Then define

$$R_\lambda = \Phi^*.$$

Note that any other sequence $\rho'_i \uparrow \lambda$ gives rise to an equivalent operation by Lemma 3.11. Also, it is easy to check that $R_\rho \geq R_{\rho'}$ if $\rho \geq \rho'$.

For each $\rho < \omega_1$, let $\mathcal{R}^\rho = \Sigma_1^{R_\rho}$ and $B\mathcal{R}^\rho = \Delta_1^{R_\rho}$. Let $\mathcal{B}\mathcal{R}^\rho$ be the least class containing clopen relations and closed under R_ρ and complementation. Thus, for instance, $\mathcal{R}^0 = \Sigma_1^1$, $B\mathcal{R}^0 = \Delta_1^1$ and $\mathcal{B}\mathcal{R}^0 = C$ -sets of Selivanovskii. Finally, set

$$\mathcal{R} = \bigcup_{\rho < \omega_1} \mathcal{R}^\rho.$$

Members of \mathcal{R} are known classically as the R -sets. It is not difficult to see that for each ρ ,

$$\mathcal{R}^\rho \subseteq \mathcal{B}\mathcal{R}^\rho \subseteq B\mathcal{R}^{\rho+1} \subseteq \mathcal{R}^{\rho+1}.$$

In fact, the inclusions can be shown to be strict (cf. [12, 13]).

The following is immediate.

LEMMA 4.1. For every ρ , R_ρ is normal and $RR_\rho \sim R_\rho$.

For each ρ , the class \mathcal{BR}^ρ can be decomposed into a hierarchy just as in the case of the Borel class. Suppose $R^\rho = \Phi^*$. We set $\mathcal{R}_0^\rho = \mathcal{R}^\rho$ and take

$$\mathcal{R}_\mu^\rho = R^\rho \left[\left(\bigcup_{\nu < \mu} \mathcal{R}_\nu^\rho \right)^c \right], \quad B\mathcal{R}_\mu^\rho = \{ E : E, E^c \in \mathcal{R}_\mu^\rho \}.$$

\mathcal{BR}_μ^ρ is the smallest class containing \mathcal{R}_μ^ρ and closed under Φ and complementation. It is immediate that these classes are included in \mathcal{BR}^ρ and indeed that $\mathcal{BR}^\rho = \bigcup_{\mu < \omega_1} \mathcal{R}_\mu^\rho$. As above we have, for each $\rho, \mu < \omega_1$,

$$\mathcal{R}_\mu^\rho \subsetneq \mathcal{BR}_\mu^\rho \subsetneq B\mathcal{R}_{\mu+1}^\rho \subsetneq \mathcal{R}_{\mu+1}^\rho.$$

We shall now show that each class $c\mathcal{R}^\rho = \Pi_1^{\mathcal{R}^\rho}$ has the pre-well-ordering property. The key to this is the following lemma.

If $E = \mathcal{R}_{\mathcal{N}}(\{E_n : n \in \omega\})$, where $\mathcal{N} = \{N_p : p \in \omega\}$ and Γ the canonical inductive operator associated with $\{E_n\}$ and \mathcal{N} , then we have $x \notin E \leftrightarrow e \in \Gamma_x^\infty$. Put

$$|(s, x)|_\Gamma = \begin{cases} \text{least } \rho \text{ such that } s \in \Gamma_x^\rho & \text{if such } \rho \text{ exists,} \\ \omega_1 & \text{otherwise.} \end{cases}$$

Thus $|(e, x)|_\Gamma < \omega_1 \leftrightarrow x \notin E$.

LEMMA 4.2 (COMPARISON OF INDICES). *Let $\mathcal{E} = \{E_n : n \in \omega\}$ and $\mathcal{F} = \{F_n : n \in \omega\}$ be two families of subsets of X and further assume that \mathcal{F} is regular, i.e., $F_t \subseteq F_s$ if $s < t$. Let $\mathcal{N} = \{N_p : p \in \omega\}$, $\mathcal{M} = \{M_p : p \in \omega\}$ be two sequences of bases. Define a sequence of bases $\{K_s : s \in \omega\}$ as follows. If $s = \langle \langle n_0, m_0 \rangle, \dots, \langle n_k, m_k \rangle \rangle$, then K_s is the canonical base for the positive analytical operation Φ defined by*

$$\Phi(\{G_n : n \in \omega\}) = \Phi_{M_{\langle m_0, \dots, m_k \rangle}}^0 \left(\left\{ \Phi_{N_{\langle n_0, \dots, n_k \rangle}}(\{G_{\langle n, m \rangle} : n \in \omega\}) : m \in \omega \right\} \right).$$

Otherwise, $K_s = \{ \omega \}$. Let

$$\mathcal{X} = \{K_s : s \in \omega\},$$

$$H_s = \begin{cases} E_{\langle n_0, \dots, n_k \rangle} \cup F_{\langle m_0, \dots, m_k \rangle}^c & \text{if } s = \langle \langle n_0, m_0 \rangle, \dots, \langle n_k, m_k \rangle \rangle, \\ X & \text{otherwise.} \end{cases}$$

Suppose Γ is the canonical inductive operator associated with \mathcal{E} and \mathcal{N} ; Δ the inductive operator associated with \mathcal{F} and \mathcal{M} . Then,

$$\{x : |(e, x)|_\Gamma < |(e, x)|_\Delta\}^c = \mathcal{R}_{\mathcal{X}}(\{H_s : s \in \omega\}).$$

PROOF. First note that

$$\eta \in K_{\langle \langle n_0, m_0 \rangle, \dots, \langle n_{i-1}, m_{i-1} \rangle \rangle}$$

$$\leftrightarrow (\exists \eta' \in M_{\langle m_0, \dots, m_{i-1} \rangle}^0)(\forall m \in \eta')(\exists \eta'' \in N_{\langle n_0, \dots, n_{i-1} \rangle})(\forall n \in \eta'')[\langle n, m \rangle \in \eta].$$

The operators Γ and Δ are as follows.

$$\begin{aligned} s \in \Gamma_x(A) &\leftrightarrow x \notin E_s \vee (\forall \eta \in N_s)(\exists n \in \eta)[s * \langle n \rangle \in A], \\ s \in \Delta_x(A) &\leftrightarrow x \notin F_s \vee (\forall \eta \in M_s)(\exists n \in \eta)[s * \langle n \rangle \in A]. \end{aligned}$$

To obtain the result look at the canonical inductive operator associated with \mathcal{K} and $\mathcal{H} = \{H_s : s \in \omega\}$:

$$s \in \Lambda_x(A) \leftrightarrow x \notin H_s \vee (\forall \eta \in K_s)(\exists n \in \eta)[s * \langle n \rangle \in A].$$

We claim that for all $x \in X$ and $n_0, m_0, \dots, n_{i-1}, m_{i-1}$

$$(8) \quad \begin{aligned} |(\langle n_0, \dots, n_{i-1} \rangle, x)|_\Gamma &< |(\langle m_0, \dots, m_{i-1} \rangle, x)|_\Delta \\ &\leftrightarrow \langle \langle n_0, m_0 \rangle, \dots, \langle n_{i-1}, m_{i-1} \rangle \rangle \in \Lambda_x^\infty \end{aligned}$$

and the result follows by taking $i = 0$.

We shall prove the implication (\rightarrow) by induction on $|(\langle n_0, \dots, n_{i-1} \rangle, x)|_\Gamma$. Suppose

$$\rho = |(\langle n_0, \dots, n_{i-1} \rangle, x)|_\Gamma < |(\langle m_0, \dots, m_{i-1} \rangle, x)|_\Delta$$

and assume, to the contrary, $\langle \langle n_0, m_0 \rangle, \dots, \langle n_{i-1}, m_{i-1} \rangle \rangle \notin \Lambda_x^\infty$. Then

$$(i) \quad x \in E_{\langle n_0, \dots, n_{i-1} \rangle} \cup F_{\langle m_0, \dots, m_{i-1} \rangle}^c \& (\exists \eta \in K_{\langle \langle n_0, m_0 \rangle, \dots, \langle n_{i-1}, m_{i-1} \rangle \rangle}) \\ (\forall s \in \eta)[\langle \langle n_0, m_0 \rangle, \dots, \langle n_{i-1}, m_{i-1} \rangle \rangle * \langle s \rangle \notin \Lambda_x^\infty].$$

Now, $\langle m_0, \dots, m_{i-1} \rangle \notin \Delta_x^\rho$ and so

$$x \in F_{\langle m_0, \dots, m_{i-1} \rangle} \& (\exists \eta' \in M_{\langle m_0, \dots, m_{i-1} \rangle})(\forall m \in \eta')[\langle m_0, \dots, m_{i-1}, m \rangle \notin \Delta_x^{\leq \rho}].$$

This implies

$$(ii) \quad x \in F_{\langle m_0, \dots, m_{i-1} \rangle} \& (\exists \eta' \in M_{\langle m_0, \dots, m_{i-1} \rangle})(\forall m \in \eta')[\rho \leq |(\langle m_0, \dots, m_{i-1}, m \rangle, x)|_\Delta].$$

Clearly from (i) and (ii), $x \in E_{\langle n_0, \dots, n_{i-1} \rangle}$ and since $\langle n_0, \dots, n_{i-1} \rangle \in \Gamma_x^\rho$,

$$(\forall \eta'' \in N_{\langle n_0, \dots, n_{i-1} \rangle})(\exists n \in \eta'')[\langle n_0, \dots, n_{i-1}, n \rangle \in \Gamma_x^{\leq \rho}]$$

which gives

$$(iii) \quad (\forall \eta'' \in N_{\langle n_0, \dots, n_{i-1} \rangle})(\exists n \in \eta'')[|(\langle n_0, \dots, n_{i-1}, n \rangle, x)|_\Gamma < \rho = |(\langle n_0, \dots, n_{i-1} \rangle, x)|_\Gamma].$$

Fix $\eta \in K_{\langle \langle n_0, m_0 \rangle, \dots, \langle n_{i-1}, m_{i-1} \rangle \rangle}$ to satisfy (i). Get $\eta' \in M_{\langle m_0, \dots, m_{i-1} \rangle}^0$ such that

$$(iv) \quad (\forall m \in \eta')(\exists \eta'' \in N_{\langle n_0, \dots, n_{i-1} \rangle})(\forall n \in \eta'')[\langle n, m \rangle \in \eta].$$

Clearly, (ii) and (1) of §2 yield $m^* \in \eta'$ such that

$$\rho \leq |(\langle m_0, \dots, m_{i-1}, m^* \rangle, x)|_\Delta.$$

By (iv) corresponding to m^* get $\eta'' \in N_{\langle n_0, \dots, n_{i-1} \rangle}$ such that $(\forall n \in \eta'')[\langle n, m^* \rangle \in \eta]$. By (iii) get $n^* \in \eta''$ such that $|(\langle n_0, \dots, n_{i-1}, n^* \rangle, x)|_\Gamma < \rho$. Clearly, $\langle n^*, m^* \rangle \in \eta$ and

$$|(\langle n_0, \dots, n_{i-1}, n^* \rangle, x)|_\Gamma < \rho \leq |(\langle m_0, \dots, m_{i-1}, m^* \rangle, x)|_\Delta$$

and by the induction hypothesis this implies

$$\langle \langle n_0, m_0 \rangle, \dots, \langle n_{i-1}, m_{i-1} \rangle, \langle n^*, m^* \rangle \rangle \in \Lambda_x^\infty.$$

This clearly contradicts our choice of η . Hence $\langle\langle n_0, m_0 \rangle, \dots, \langle n_{i-1}, m_{i-1} \rangle\rangle \in \Lambda_x^\infty$.

To prove the other implication, set

$$\begin{aligned} A^* = \{ \langle\langle n_0, m_0 \rangle, \dots, \langle n_{i-1}, m_{i-1} \rangle\rangle : \\ |(\langle n_i, \dots, n_{i-1} \rangle, x)|_\Gamma < |(\langle m_0, \dots, m_{i-1} \rangle, x)|_\Delta \} \\ \cup \{ t : t \text{ is not of the form } \langle\langle n_0, m_0 \rangle, \dots, \langle m_{i-1}, m_{i-1} \rangle\rangle \}. \end{aligned}$$

We shall show that $\Lambda_x(A^*) \subseteq A^*$, from which it will follow that $\Lambda_x^\infty \subseteq A^*$.

So let

$$(v) \quad \langle\langle n_0, m_0 \rangle, \dots, \langle n_{i-1}, m_{i-1} \rangle\rangle \in \Lambda_x(A^*).$$

We will have to show that

$$|(\langle n_0, \dots, n_{i-1} \rangle, x)|_\Gamma < |(\langle m_0, \dots, m_{i-1} \rangle, x)|_\Delta.$$

Assume to the contrary that

$$|(\langle m_0, \dots, m_{i-1} \rangle, x)|_\Delta \leq |(\langle n_0, \dots, n_{i-1} \rangle, x)|_\Gamma.$$

From (v) we have

$$\begin{aligned} x \in E_{\langle n_0, \dots, n_{i-1} \rangle}^c \cap F_{\langle m_0, \dots, m_{i-1} \rangle} \vee (\forall \eta \in K_{\langle\langle n_0, m_0 \rangle, \dots, \langle n_{i-1}, m_{i-1} \rangle\rangle})(\exists s \in \eta) \\ [(\langle n_0, m_0 \rangle, \dots, \langle n_{i-1}, m_{i-1} \rangle, s) \in A^*]. \end{aligned}$$

If $x \in E_{\langle n_0, \dots, n_{i-1} \rangle}^c \cap F_{\langle m_0, \dots, m_{i-1} \rangle}$, then $|(\langle n_0, \dots, n_{i-1} \rangle, x)|_\Gamma = 0$ and

$$|(\langle m_0, \dots, m_{i-1} \rangle, x)|_\Delta > 0$$

and we are done. So assume

$$\begin{aligned} (vi) \quad x \in (E_{\langle n_0, \dots, n_{i-1} \rangle} \cup F_{\langle m_0, \dots, m_{i-1} \rangle}^c) \& (\forall \eta \in K_{\langle\langle n_0, m_0 \rangle, \dots, \langle n_{i-1}, m_{i-1} \rangle\rangle})(\exists s \in \eta) \\ [(\langle n_0, m_0 \rangle, \dots, \langle n_{i-1}, m_{i-1} \rangle, s) \in A^*]. \end{aligned}$$

If $x \in F_{\langle m_0, \dots, m_{i-1} \rangle}^c$, then by regularity $x \in F_{\langle m_0, \dots, m_{i-1}, m \rangle}^c$ for all m , and hence $(\forall m)[|(\langle m_0, \dots, m_{i-1}, m \rangle, x)|_\Delta = 0]$. But this is not possible by (vi). Therefore,

$$\begin{aligned} (vii) \quad x \in (E_{\langle n_0, \dots, n_{i-1} \rangle} \cap F_{\langle m_0, \dots, m_{i-1} \rangle}) \\ \& (\forall \eta \in K_{\langle\langle n_0, m_0 \rangle, \dots, \langle n_{i-1}, m_{i-1} \rangle\rangle})(\exists s \in \eta) \\ [(\langle n_0, m_0 \rangle, \dots, \langle n_{i-1}, m_{i-1} \rangle, s) \in A^*]. \end{aligned}$$

Case 1. $|(\langle n_0, \dots, n_{i-1} \rangle, x)|_\Gamma = \omega_1$. In this case $\langle n_0, \dots, n_{i-1} \rangle \notin \Gamma_x^\infty$ and hence

$$(\exists \eta'' \in N_{\langle n_0, \dots, n_{i-1} \rangle})(\forall n \in \eta'')[|(\langle n_0, \dots, n_{i-1}, n \rangle, x)|_\Gamma = \omega_1].$$

Fix such an $\eta'' \in N_{\langle n_0, \dots, n_{i-1} \rangle}$. Pick any $\eta' \in M_{\langle m_0, \dots, m_{i-1} \rangle}^0$ and put

$$\eta^* = \{ \langle n, m \rangle : n \in \eta'' \& m \in \eta' \}.$$

Clearly $\eta^* \in K_{\langle \langle n_0, m_0 \rangle, \dots, \langle n_{i-1}, m_{i-1} \rangle \rangle}$ and, moreover,

$$(\forall \langle n, m \rangle \in \eta^*) [|(\langle m_0, \dots, m_{i-1}, m \rangle, x)|_\Delta \leq \omega_1 = |(\langle n_0, \dots, n_{i-1}, n \rangle, x)|_\Gamma].$$

This contradicts (vii).

Case 2. $|(\langle n_0, \dots, n_{i-1} \rangle, x)|_\Gamma = \rho < \omega_1$. Here we have, by our assumption, $\langle m_0, \dots, m_{i-1} \rangle \in \Delta_x^\rho$ and since $x \in F_{\langle m_0, \dots, m_{i-1} \rangle}$,

$$(\forall \eta' \in M_{\langle m_0, \dots, m_{i-1} \rangle}) (\exists m \in \eta') [|(\langle m_0, \dots, m_{i-1}, m \rangle, x)|_\Delta < \rho]$$

which implies by (2) that

$$\begin{aligned} & (\exists \eta' \in M_{\langle m_0, \dots, m_{i-1} \rangle}^0) (\forall m \in \eta') [|(\langle m_0, \dots, m_{i-1}, m \rangle, x)|_\Delta < \rho] \text{ hence,} \\ & (\exists \eta' \in M_{\langle m_0, \dots, m_{i-1} \rangle}^0) (\forall m \in \eta') [\langle n_0, \dots, n_{i-1} \rangle \notin \Gamma_x^{(|(\langle m_0, \dots, m_{i-1}, m \rangle, x)|_\Delta)}]. \end{aligned}$$

Consequently,

$$\begin{aligned} & (\exists \eta' \in M_{\langle m_0, \dots, m_{i-1} \rangle}^0) (\forall m \in \eta') (\exists \eta'' \in N_{\langle n_0, \dots, n_{i-1} \rangle}) (\forall n \in \eta'') \\ & [|(\langle m_0, \dots, m_{i-1}, m \rangle, x)|_\Delta \leq |(\langle n_0, \dots, n_{i-1}, n \rangle, x)|_\Gamma]. \end{aligned}$$

This clearly contradicts (vii). Thus in either case we have a contradiction and so $\Lambda_x^\infty(A^*) \subseteq A^*$. Thus $\Lambda_x^\infty \subseteq A^*$. This proves the other implication of (8).

The following trick is due to Lyapunov.

LEMMA 4.3 (INCREASING THE INDEX BY 1). *Let Γ be the canonical inductive operator associated with $\mathcal{E} = \{ E_p : p \in \omega \}$ and $\mathcal{N} = \{ N_k : k \in \omega \}$. Define*

$$\begin{aligned} E_{\langle \rangle}^* &= X; \\ E_s^* &= \begin{cases} E_{\langle n_1, \dots, n_k \rangle} & \text{if } s = \langle 1, n_1, \dots, n_k \rangle, \\ \emptyset & \text{otherwise.} \end{cases} \end{aligned}$$

Put

$$\begin{aligned} N_{\langle \rangle}^* &= \{ \{ 1 \} \}; \\ N_s^* &= \begin{cases} N_{\langle n_1, \dots, n_k \rangle} & \text{if } s = \langle 1, n_1, \dots, n_k \rangle, \\ \{ \{ n \} : n \in \omega \} & \text{otherwise.} \end{cases} \end{aligned}$$

If $\mathcal{N}^* = \{ N_k^* : k \in \omega \}$, then

$$\mathcal{R}_{\mathcal{N}^*}(\{ E_s^* : s \in \omega \}) = \mathcal{R}_{\mathcal{N}}(\{ E_s : s \in \omega \})$$

and $|(\langle e, x \rangle)_{\Gamma^*} = |(\langle e, x \rangle)_\Gamma + 1$, where Γ^* is the inductive operator associated with

$$\mathcal{E}^* = \{ E_s^* : s \in \omega \} \text{ and } \mathcal{N}^*.$$

PROOF. We shall prove by induction on ρ that

$$\langle 1, n_1, \dots, n_k \rangle \in \Gamma_x^{*\rho} \leftrightarrow \langle n_1, \dots, n_k \rangle \in \Gamma_x^\rho.$$

Now,

$$\begin{aligned}
\langle 1, n_1, \dots, n_k \rangle \in \Gamma_x^{*\rho} &\leftrightarrow x \notin E_{\langle 1, n_1, \dots, n_k \rangle}^* \vee \\
&\quad (\forall \eta \in N_{\langle 1, n_1, \dots, n_k \rangle}^*)(\exists n \in \eta)[\langle 1, n_1, \dots, n_k, n \rangle \in \Gamma_x^{* < \rho}] \\
&\leftrightarrow x \notin E_{\langle n_1, \dots, n_k \rangle} \vee \\
&\quad (\forall \eta \in N_{\langle n_1, \dots, n_k \rangle})(\exists n \in \eta)[\langle n_1, \dots, n_k, n \rangle \in \Gamma_x^{* < \rho}] \\
&\leftrightarrow \langle n_1, \dots, n_k \rangle \in \Gamma_x^\rho. \\
\therefore e \in \Gamma_x^{*\rho+1} &\leftrightarrow (\forall \eta \in N_e^*)(\exists n \in \eta)[\langle n \rangle \in \Gamma_x^{*\rho}] \\
&\leftrightarrow \langle 1 \rangle \in \Gamma_x^{*\rho} \\
&\leftrightarrow e \in \Gamma_x^\rho.
\end{aligned}$$

Hence $\mathcal{R}_{\mathcal{N}^*}(\{E_s^* : s \in \omega\}) = \mathcal{R}_{\mathcal{N}}(\{E_s : s \in \omega\})$ and $\|e, x\|_{\Gamma^*} = \|e, x\|_{\Gamma} + 1$.

The following lemma follows from above. One has only to observe that for any positive analytical operation Φ , $R\Phi$ is normal, $R\Phi^0\Phi \sim R\Phi\Phi^0$ and if

$$\mathcal{N} = \{N_s : s \in \omega\}$$

is a family of bases such that $\Phi \geq \Phi_{N_s}$ for each s , then $R\Phi \geq \mathcal{R}_{\mathcal{N}}$.

LEMMA 4.4. *Let $\Phi \geq \cup$ be a positive analytical operation and $\mathcal{N} = \{N_s : s \in \omega\}$, $\mathcal{M} = \{M_s : s \in \omega\}$ be two families of bases such that for each s , Φ^* subsumes Φ_{N_s} and Φ_{M_s} . Suppose $\{E_s : s \in \omega\}$ is a family of sets in $\Sigma_1^{\Phi^*}$ and $\{F_s : s \in \omega\}$ a regular family in $\Pi_1^{\Phi^*}$. Let Γ be the canonical inductive operator associated with $\mathcal{E} = \{E_s : s \in \omega\}$ and \mathcal{N} , and Δ that associated with $\mathcal{F} = \{F_s : s \in \omega\}$ and \mathcal{M} . Put $\beta_1(x) = \|e, x\|_{\Gamma}$ and $\beta_2(x) = \|e, x\|_{\Delta}$. Then*

- (a) $\{x : \beta_1(x) < \beta_2(x)\} \in \Pi_1^{\Phi^*}$,
- (b) $\{x : \beta_1(x) < \omega_1 \ \& \ \beta_1(x) \leq \beta_2(x)\} \in \Pi_1^{\Phi^*}$.

PROOF. The first assertion follows from the comparison of indices lemma and the observations made above. The second assertion follows from the first by increasing the index of Δ by 1.

By slightly modifying the inductive operator Λ in the proof of Lemma 4.2, one can obtain the following

COROLLARY 4.5. *Let Γ, Δ, Φ be as in 4.4. Then*

$$\{(s, t, x) : \|s, x\|_{\Gamma} < \|t, x\|_{\Delta}\} \in \Pi_1^{\Phi^*}.$$

THEOREM 4.6. *For any positive analytical operation $\Phi \geq \cup$, $\Pi_1^{\Phi^*}$ has the pre-well-ordering property.*

PROOF. Let $E \in \Pi_1^{\Phi^*}$ and suppose

$$E^c = \Phi^*(\{A_s : s \in \omega\}) \quad \text{with } \{A_s\} \text{ clopen and regular.}$$

Let $\beta(x)$ be the norm on E induced by the canonical inductive operator. Let N be the canonical base for Φ and put $\mathcal{N} = \{NN^0\}$, $\mathcal{M} = \{NN^0\}$. For each s , set

$$E_s = A_s \times X, \quad F_s = X \times A_s$$

and let $\beta_1(x, y), \beta_2(x, y)$ be the norms induced by the inductive operators associated with $\{E_s : s \in \omega\}, \mathcal{N}$ and $\{F_s : s \in \omega\}, \mathcal{M}$. Since all the hypotheses of Lemma 4.4 are satisfied, the sets $\{(x, y) : \beta_1(x, y) < \beta_2(x, y)\}$ and

$$\{(x, y) : \beta_1(x, y) < \omega_1 \& \beta_1(x, y) \leq \beta_2(x, y)\}$$

are in $\Pi_1^{\Phi^*}$. But

$$\{(x, y) : \beta_1(x, y) < \beta_2(x, y)\} = \{(x, y) : \beta(x) < \beta(y)\}$$

and

$$\begin{aligned} \{(x, y) : \beta_1(x, y) < \omega_1 \& \beta_1(x, y) \leq \beta_2(x, y)\} \\ = \{(x, y) : \beta(x) < \omega_1 \& \beta(x) \leq \beta(y)\}. \end{aligned}$$

Consequently, $\Pi_1^{\Phi^*}$ is normed.

COROLLARY 4.7. *For each $\rho < \omega_1$, $c\mathcal{R}^\rho$ has the pre-well-ordering property.*

5. Complexity of winning strategies and the transfer property of the R-operator.

THEOREM 5.1. *Let Φ be a positive analytical operation which subsumes both (countable) \cup and \cap . Let ∇ be the σ -field generated by $\Sigma_1^{\Phi^*}$. Let $E \in \Sigma_1^{\Phi^*}$ be such that*

$$(i) \quad x \in E \leftrightarrow (\exists \eta_0 \in NN^0)(\forall n_0 \in \eta_0)(\exists \eta_1 \in NN^0)(\forall n_1 \in \eta_1) \cdots \cdots (\forall k) [x \in E_{\langle n_0, \dots, n_{k-1} \rangle}],$$

N being the canonical base for Φ . Then, there is a ∇ -measurable function $x \mapsto \sigma_x$ such that σ_x is a winning strategy for the player \exists , whenever $x \in E$.

PROOF. Let Γ be the canonical inductive operator associated with NN^0 and $\{E_s : s \in \omega\}$ and put $NN^0 = M$. Define

$$\beta(s, x) = \begin{cases} \text{least } \rho \text{ such that } s \in \Gamma_x^\rho & \text{if } s \in \Gamma_x^\infty, \\ \omega_1 & \text{otherwise.} \end{cases}$$

Now suppose $x \in E$. So \exists wins the game (i).

If $\eta_0, \eta_1, \dots, \eta_{k-1}$ and n_0, n_1, \dots, n_{k-1} are the first k relevant moves of \exists and \forall , notice that \exists goes on to win the game (i), i.e., he is in a winning position iff $\langle n_0, \dots, n_{k-1} \rangle \notin \Gamma_x^\infty$ i.e., iff $\beta(\langle n_0, \dots, n_{k-1} \rangle, x) = \omega_1$. In such a case, \exists has to play an $\eta \in M$ such that $(\forall n \in \eta)[\beta(\langle n_0, \dots, n_{k-1}, n \rangle, x) = \omega_1]$. We, therefore, define for each x , the strategy σ_x for \exists as follows:

$$p \in \sigma_x(s) \leftrightarrow \beta(s, x) \leq \beta(s * \langle p \rangle, x).$$

Clearly by 4.5, the map $x \mapsto \sigma_x$ is ∇ -measurable. We shall now show that if $x \in E$, then σ_x is a winning strategy for \exists in the game (i). Suppose $\eta_0, n_0, \eta_1, n_1, \dots, \eta_{k-1}, n_{k-1}$ are the first k moves of \exists and \forall and assume that \exists has not yet lost the game i.e. he is in a winning position. Consequently, we have $(\beta(\langle n_0, n_1, \dots, n_{k-1} \rangle, x) = \omega_1$ and hence $\langle n_0, \dots, n_{k-1} \rangle \notin \Gamma_x^\infty$. Therefore,

$$(ii) \quad (\exists \eta \in M)(\forall n \in \eta)[\beta(\langle n_0, n_1, \dots, n_{k-1}, n \rangle, x) = \omega_1].$$

By definition,

$$\begin{aligned}\sigma_x(\langle n_0, \dots, n_{k-1} \rangle) &= \{ p : \beta(\langle n_0, \dots, n_{k-1} \rangle, x) \leq \beta(\langle n_0, \dots, n_{k-1}, p \rangle, x) \} \\ &= \{ p : \beta(\langle n_0, \dots, n_{k-1}, p \rangle, x) = \omega_1 \}.\end{aligned}$$

Hence, by (ii) and the completeness of M ,

$$\sigma_x(\langle n_0, \dots, n_{k-1} \rangle) = \eta \in M,$$

and moreover, $\forall p \in \eta$, $\beta(\langle n_0, \dots, n_{k-1}, p \rangle, x) = \omega_1$, so that \exists is still in a winning position.

REMARK. Notice that

(iii)

$$\begin{aligned}x \notin E &\leftrightarrow (\forall \eta_0 \in M)(\exists n_0 \in \eta_0)(\forall \eta_1 \in M)(\exists n_1 \in \eta_1) \cdots (\exists k) \left[x \notin E_{\langle n_0, \dots, n_{k-1} \rangle} \right] \\ &\leftrightarrow (\exists \eta_0 \in M^0)(\forall n_0 \in \eta_0) \cdots (\exists k) \left[x \notin E_{\langle n_0, \dots, n_{k-1} \rangle} \right].\end{aligned}$$

Here also we can have a definable winning strategy for \exists whenever $x \notin E$. Unlike the game (i), here \exists has to play such that at each stage the value of β is *decreased*. The following will give a ∇ -measurable winning strategy for \exists in the game (iii) whenever $x \notin E$:

$$p \in \sigma_x(s) \leftrightarrow x \notin E_s \vee (\beta(s * \langle p \rangle, x) < \beta(s, x)).$$

DEFINITION 5.2. A set $E \subseteq X \times Y$ is said to be normal if for each $x \in X$, $E^x = \{ y : (x, y) \in E \}$ has the Baire property. If E is normal and $U \subseteq Y$ is open, then define

$$E^{*U} = \{ x \in X : E^x \text{ is comeager in } U \}.$$

If $U = Y$, we write E^* instead of E^{*U} .

LEMMA 5.3. Let Φ be a positive analytical operation which preserves the Baire property and let $E = \Phi^*(\{ E_s : s \in \omega \})$, with each $E_s \subseteq \omega^\omega \times \omega^\omega$ normal. Define E_s^μ and T^μ as in 3.6. Then for any $s \in \omega$,

$$E * \Sigma(s) = \bigcap_{\mu < \omega_1} [E_s^\mu]^{*\Sigma(s)} = \bigcup_{\mu < \omega_1} [E_s^\mu - T^\mu]^{*\Sigma(s)}.$$

PROOF. As in Theorem 3.8, one can easily check that E_s^μ and T^μ are normal for each μ . Since $E = \bigcap_{\mu < \omega_1} E_s^\mu$, it follows that

$$(i) \quad E * \Sigma(s) \subseteq \bigcap_{\mu < \omega_1} [E_s^\mu]^{*\Sigma(s)}.$$

Next, suppose $x \in [E_s^\mu]^{*\Sigma(s)}$ for all $\mu < \omega_1$. For each $p \in \omega$, $\langle (E_p^\mu)^x : \mu < \omega_1 \rangle$ is a decreasing sequence of sets with Baire property. Hence by the countable chain condition, $\exists \beta(p) < \omega_1$ such that

$$(\forall \rho > \beta(p)) \left[(E_p^{\beta(p)} - E_p^\rho)^x \text{ is meager} \right].$$

Choose ρ_0 such that $\beta(p) < \rho_0 < \omega_1$, for all p . Then $(\forall p) [(E_p^{\rho_0} - E_p^{\rho_0+1})^x \text{ is meager}]$ and hence $(T^{\rho_0})^x$ is meager. Since $x \in [E_s^{\rho_0}]^{*\Sigma(s)}$, $(E_s^{\rho_0})^x$ is comeager in $\Sigma(s)$. Therefore, $(E_s^{\rho_0})^x - (T^{\rho_0})^x$ is comeager in $\Sigma(s)$ and so $x \in [E_s^{\rho_0} - T^{\rho_0}]^{*\Sigma(s)}$.

Thus,

$$(ii) \quad \bigcap_{\mu < \omega_1} [E_\epsilon^\mu]^{\star \Sigma(s)} \subseteq \bigcup_{\mu < \omega_1} [E_\epsilon^\mu - T^\mu]^{\star \Sigma(s)}.$$

Finally, since $(E_\epsilon^\mu - T^\mu) \subseteq E$ for each μ ,

$$(iii) \quad \bigcup_{\mu < \omega_1} [E_\epsilon^\mu - T^\mu]^{\star \Sigma(s)} \subseteq E^{\star \Sigma(s)}.$$

The result now follows from (i)–(iii).

TRANSFER THEOREM 5.4. *Let Φ and Ψ be two positive analytical operations such that Φ preserves the Baire property and Ψ is normal and subsumes both $(\text{countable}) \cup$ and \cap . Suppose, moreover, that there are functions f and g such that for any normal family $\{E_p : p \in \omega\}$ of subsets of $\omega^\omega \times \omega^\omega$ with $E = \Phi(\{E_p : p \in \omega\})$,*

$$E^{\star \Sigma(s)} = \Psi\left(\left\{E_{f(p)}^{\star \Sigma(\hat{s}g(p))} : p \in \omega\right\}\right).$$

Then for any normal family $\{F_p : p \in \omega\}$ of subsets of $\omega^\omega \times \omega^\omega$,

(a) $F = \Phi^0(\{F_p : p \in \omega\})$ *implies that*

$$F^{\star \Sigma(s)} = \Psi^0\left(\left\{F_{\gamma(p)}^{\star \Sigma(\hat{s}\delta(p))} : p \in \omega\right\}\right),$$

for suitable functions γ and δ (independent of the family $\{F_p\}$).

(b) $F = \Phi \cdot \Phi^0(\{F_p : p \in \omega\})$ *implies*

$$F^{\star \Sigma(s)} = \Psi \Psi^0\left(\left\{F_{\alpha(p)}^{\star \Sigma(\hat{s}\beta(p))} : p \in \omega\right\}\right),$$

where

$$\alpha(s) = \begin{cases} \langle f(n), \gamma(m) \rangle & \text{if } s = \langle n, m \rangle, \\ \text{arbitrary} & \text{otherwise;} \end{cases}$$

$$\beta(s) = \begin{cases} g(n) \hat{\delta}(m) & \text{if } s = \langle n, m \rangle, \\ \text{arbitrary} & \text{otherwise.} \end{cases}$$

(c) $F = \Phi^*(\{F_p : p \in \omega\})$ *implies*

$$F^{\star \Sigma(s)} = \Psi^*\left(\left\{F_{f(p)}^{\star \Sigma(\hat{s}\tilde{g}(p))} : p \in \omega\right\}\right),$$

where

$$\tilde{f}(s) = \begin{cases} \langle \alpha(n_0), \dots, \alpha(n_{k-1}) \rangle & \text{if } s = \langle n_0, \dots, n_{k-1} \rangle, \\ \text{arbitrary} & \text{otherwise;} \end{cases}$$

$$\tilde{g}(s) = \begin{cases} \beta(n_0) \hat{\beta}(n_1) \cdots \hat{\beta}(n_{k-1}) & \text{if } s = \langle n_0, \dots, n_{k-1} \rangle, \\ \text{arbitrary} & \text{otherwise.} \end{cases}$$

PROOF. Let N and M be the canonical bases for Φ and Ψ , respectively.

(a) Set $G = F^c$ and $G_p = F_p^c$ for each p .

Then $G = \Phi(\{G_p : p \in \omega\})$; and since Φ preserves the Baire property, each G^x has

the Baire property. Therefore,

G^x is nonmeager in $\Sigma(s)$

$$\begin{aligned} &\leftrightarrow (\exists u) [G^x \text{ is comeager in } \Sigma(\hat{s}u)] \\ &\leftrightarrow (\exists u)(\exists \eta \in M)(\forall m \in \eta) [G_{f(m)}^x \text{ is comeager in } \Sigma(\hat{s}\hat{u}g(m))] \quad \text{by hypothesis} \\ &\leftrightarrow (\exists u)(\exists \eta \in M)(\forall m \in \eta)(\forall v) [G_{f(m)}^x \text{ is nonmeager in } \Sigma(\hat{s}\hat{u}g(m)\hat{v})] \\ &\leftrightarrow (\exists \eta \in M)(\forall m \in \eta) [G_{\gamma(m)}^x \text{ is nonmeager in } \Sigma(\hat{s}\hat{\delta}(m))], \end{aligned}$$

for some functions γ and δ , as $\Phi_M \geq \cap, \cup$ and is normal. Hence,

$$\begin{aligned} F^x \text{ is comeager in } \Sigma(s) &\leftrightarrow G^x \text{ is meager in } \Sigma(s) \\ &\leftrightarrow (\forall \eta \in M)(\exists m \in \eta) [G_{\gamma(m)}^x \text{ is meager in } \Sigma(\hat{s}\hat{\delta}(m))] \\ &\leftrightarrow (\forall \eta \in M)(\exists m \in \eta) [F_{\gamma(m)}^x \text{ is comeager in } \Sigma(\hat{s}\hat{\delta}(m))] \\ &\leftrightarrow (\exists \eta \in M^0)(\forall m \in \eta) [F_{\gamma(m)}^x \text{ is comeager in } \Sigma(\hat{s}\hat{\delta}(m))]. \end{aligned}$$

This proves (a).

To prove (b), observe that

F^x is comeager in $\Sigma(s)$

$$\begin{aligned} &\leftrightarrow [\Phi(\{ \Phi^0(\{ F_{\langle p, q \rangle} : q \in \omega \}) : p \in \omega \})]^x \text{ is comeager in } \Sigma(s) \\ &\leftrightarrow (\exists \eta \in M)(\forall n \in \eta) [(\Phi^0(\{ F_{\langle f(n), q \rangle} : q \in \omega \}))^x \text{ is comeager in } \Sigma(\hat{s}g(n))] \\ &\quad \text{(by hypothesis),} \\ &\leftrightarrow (\exists \eta \in M)(\forall n \in \eta)(\exists \xi \in M^0)(\forall m \in \xi) \\ &\quad [F_{\langle f(n), \gamma(m) \rangle}^x \text{ is comeager in } \Sigma(\hat{s}g(n)\hat{\delta}(m))] \quad \text{(by (a)).} \end{aligned}$$

Setting $\alpha(\langle n, m \rangle) = \langle f(n), \gamma(m) \rangle$, $\beta(\langle n, m \rangle) = g(n)\hat{\delta}(m)$, we get (b).

(c) Since $F = \Phi^*(\{ F_p : p \in \omega \})$, the canonical inductive operator Γ is given by

$$\begin{aligned} p \in \Gamma_z(A) &\leftrightarrow z \notin F_p \vee (\forall \eta \in N)(\exists n \in \eta)(\exists \xi \in N)(\forall m \in \xi) \\ &\quad [p * \langle \langle n, m \rangle \rangle \in A]; z \in \omega^\omega \times \omega^\omega. \end{aligned}$$

Define $z \in F_p^\mu \leftrightarrow p \notin \Gamma_z^\mu$. Then by Lemma 5.3,

$$(i) \quad F^{\bullet \Sigma(s)} = \bigcap_{\mu < \omega_1} [F_e^\mu]^{\bullet \Sigma(s)}.$$

Define a set relation Δ operative on ω as follows:

$$\begin{aligned} t \in \Delta_x(A) &\leftrightarrow F_{(t)_0}^x \text{ is not comeager in } \Sigma((t)_1) \vee \\ &\quad (\forall \eta \in M)(\exists n \in \eta)(\exists \xi \in M)(\forall m \in \xi) \\ &\quad [\langle (t)_0 \hat{\langle \langle f(n), \gamma(m) \rangle \rangle}, (t)_1 \hat{g(n)\delta(m)} \rangle \in A]. \end{aligned}$$

We shall show that

$$(ii) \quad x \notin F^{\bullet \Sigma(s)} \leftrightarrow \langle e, s \rangle \in \Delta_x^\infty.$$

To prove (ii) we shall show by induction on μ that

$$x \notin [F_t^\mu]^{\star \Sigma(s)} \leftrightarrow \langle t, s \rangle \in \Delta_x^\mu.$$

This is clearly true for $\mu = 0$, so assume $\mu > 0$. Now,

$$x \notin [F_t^\mu]^{\star \Sigma(s)} \leftrightarrow (F_t^\mu)^x \text{ is not comeager in } \Sigma(s)$$

$$\leftrightarrow F_t^x \text{ is not comeager in } \Sigma(s) \vee \left[\Phi_{NN^0} \left(\left\{ \bigcap_{\lambda < \mu} F_{t' \langle p \rangle}^\lambda : p \in \omega \right\} \right) \right]^x$$

is not comeager in $\Sigma(s)$

$$\leftrightarrow F_t^x \text{ is not comeager in } \Sigma(s) \vee (\forall \eta \in M)(\exists n \in \eta)(\exists \xi \in M)(\forall m \in \xi)$$

$$\left[\left(\bigcap_{\lambda < \mu} F_{t' \langle \langle f(n), \gamma(m) \rangle \rangle}^\lambda \right)^x \text{ is not comeager in } \Sigma(\hat{s}g(n)\hat{\delta}(m)) \right] \quad (\text{by (b)})$$

$$\leftrightarrow F_t^x \text{ is not comeager in } \Sigma(s) \vee (\forall \eta \in M)(\exists n \in \eta)(\exists \xi \in M)(\forall m \in \xi)$$

$$(\exists \lambda < \mu) \left[\left(F_{t' \langle \langle f(n), \gamma(m) \rangle \rangle}^\lambda \right)^x \text{ is not comeager in } \Sigma(\hat{s}g(n)\hat{\delta}(m)) \right]$$

$$\leftrightarrow F_t^x \text{ is not comeager in } \Sigma(s) \vee (\forall \eta \in M)(\exists n \in \eta)(\exists \xi \in M)(\forall m \in \xi)$$

$$(\exists \lambda < \mu) \left[\hat{t} \langle \langle f(n), \gamma(m) \rangle \rangle, \hat{s}g(n)\hat{\delta}(m) \rangle \in \Delta_x^\lambda \right]$$

by the induction hypothesis

$$\leftrightarrow \langle t, s \rangle \in \Delta_s^\mu.$$

Hence, putting $t = \langle \rangle$ and using (i), we obtain (ii). Therefore,

$$F^x \text{ is comeager in } \Sigma(s) \leftrightarrow x \in F^{\star \Sigma(s)} \leftrightarrow \langle e, s \rangle \notin \Delta_x^\infty \quad \text{by (ii)}$$

$$\leftrightarrow (\exists \eta_0 \in M)(\forall n_0 \in \eta_0)(\forall \xi_0 \in M)(\exists m_0 \in \xi_0) \dots$$

$$(\forall k) \left[F^x \langle \langle f(n_0), \gamma(m_0) \rangle, \dots, \langle f(n_{k-1}), \gamma(m_{k-1}) \rangle \rangle \text{ is comeager in } \Sigma(\hat{s}g(n_0)\hat{\delta}(m_0) \cdots \hat{s}g(n_{k-1})\hat{\delta}(m_{k-1})) \right].$$

Define \tilde{f} and \tilde{g} as follows:

$$\tilde{f}(s) = \begin{cases} \langle \alpha(n_0), \dots, \alpha(n_{k-1}) \rangle & \text{if } s = \langle n_0, \dots, n_{k-1} \rangle, \\ \text{arbitrary} & \text{otherwise;} \end{cases}$$

$$\tilde{g}(s) = \begin{cases} \beta(n_0)\hat{\beta}(n_1) \cdots \hat{\beta}(n_{k-1}) & \text{if } s = \langle n_0, \dots, n_{k-1} \rangle, \\ \text{arbitrary} & \text{otherwise.} \end{cases}$$

The result now follows immediately.

REMARK 5.4.1. The import of the Transfer Theorem is that if a set E computed by Φ is such that $E^{\star \Sigma(s)}$ is computed by Ψ , then for any set F computed by Φ^* , $F^{\star \Sigma(s)}$ is computed by Ψ^*

REMARK 5.4.2. It is clear from the proof of 5.4(c) that, under the hypothesis of the theorem, for any normal family $\{F_p\}$ of subsets of $\omega^\omega \times \omega^\omega$ with $F = R\Phi_N(\{F_p\})$,

$$F^{\star \Sigma(s)} = R\Phi_M \left(\left\{ F_{f' \langle p \rangle}^{\star \Sigma(\hat{s}g'(p))} : p \in \omega \right\} \right),$$

where

$$\begin{aligned} f'(\langle n_0, \dots, n_{k-1} \rangle) &= \langle f(n_0), \dots, f(n_{k-1}) \rangle, \\ g'(\langle n_0, \dots, n_{k-1} \rangle) &= g(n_0) \hat{g}(n_1) \cdots \hat{g}(n_{k-1}). \end{aligned}$$

Here, normality of $\Psi = \Phi_M$ is not required.

6. Applications. We shall now apply the Transfer Theorem to deduce Vaught's Formula for E^* (cf. [19, Theorem 1.6]) and the Category Formula of Burgess [3].

THEOREM 6.1 (VAUGHT). *Assume $E = \mathcal{A}(\{E_n : n \in \omega\})$, with each $E_n \subset \omega^\omega \times \omega^\omega$ normal. Then $x \in E^{*\Sigma(s)}$ if and only if*

$$\begin{aligned} &(\forall u_0 \in \text{Seq})(\exists v_0 \in \text{Seq})(\exists k_0)(\forall u_1 \in \text{Seq})(\exists v_1 \in \text{Seq})(\exists k_1) \cdots \\ &(\forall i) \left[E_{\langle k_0, \dots, k_{i-1} \rangle}^x \text{ is comeager in } \Sigma(\hat{s} \hat{u}_0 \hat{v}_0 \cdots \hat{u}_{i-1} \hat{v}_{i-1}) \right]. \end{aligned}$$

PROOF. Let $N = \{\langle n \rangle : n \in \omega\}$ so that $\Phi_N = \cup$. Suppose $F = \Phi_N(\{F_n : n \in \omega\}) = \cup_{n \in \omega} F_n$, where each F_n is a normal subset of $\omega^\omega \times \omega^\omega$. Then,

$$\begin{aligned} F^x \text{ is comeager in } \Sigma(s) &\leftrightarrow (\forall u \in \text{Seq}) [F^x \text{ is nonmeager in } \Sigma(\hat{s}u)] \\ &\leftrightarrow (\forall u \in \text{Seq})(\exists k) [F_k^x \text{ is nonmeager in } \Sigma(\hat{s}u)] \\ &\leftrightarrow (\forall u \in \text{Seq})(\exists k)(\exists v \in \text{Seq}) [F_k^x \text{ is comeager in } \Sigma(\hat{s}u\hat{v})]. \end{aligned}$$

Let Φ_M be a $\delta - s$ operation such that for any family $\{A_n\}$,

$$\Phi_M(\{A_n : n \in \omega\}) = \bigcap_{u \in \text{Seq}} \bigcup_{v \in \text{Seq}} \bigcup_{k \in \omega} A_{\langle u, \langle k, v \rangle \rangle}.$$

Hence,

$$\begin{aligned} &F^x \text{ is comeager in } \Sigma(s) \\ &\leftrightarrow (\exists \eta \in M)(\forall n \in \eta) [F_{(n)_{1,0}}^x \text{ is comeager in } \Sigma(\hat{s}(n)_0 \hat{(n)}_{1,1})] \\ &\leftrightarrow (\exists \eta \in M)(\forall n \in \eta) [F_{f(n)}^x \text{ is comeager in } \Sigma(\hat{s}g(n))], \end{aligned}$$

where $f(n) = (n)_{1,0}$ and $g(n) = (n)_0 \hat{(n)}_{1,1}$. Since

$$E = \mathcal{A}(\{E_n : n \in \omega\}) = R\Phi_N(\{E_n\}),$$

the Transfer Theorem immediately gives

$$E^{*\Sigma(s)} = R\Phi_M \left(\left\{ E_{\tilde{f}(p)}^{*\Sigma(s)\tilde{g}(p)} : p \in \omega \right\} \right),$$

where

$$\tilde{g}(\langle n_0, \dots, n_{k-1} \rangle) = g(n_0) \hat{g}(n_1) \cdots \hat{g}(n_{k-1})$$

and $\tilde{f}(\langle n_0, \dots, n_{k-1} \rangle) = \langle f(n_0), \dots, f(n_{k-1}) \rangle$. Therefore,

$$\begin{aligned} &E^x \text{ is comeager in } \Sigma(s) \\ &\leftrightarrow (\exists \eta_0 \in M)(\forall n_0 \in \eta_0)(\exists \eta_1 \in M)(\forall n_1 \in \eta_1) \cdots \\ &(\forall i) \left[E_{\langle f(n_0), \dots, f(n_{i-1}) \rangle}^x \text{ is comeager in } \Sigma(\hat{s}g(n_0) \hat{\cdots} \hat{g}(n_{i-1})) \right] \\ &(\forall u_0 \in \text{Seq})(\exists v_0 \in \text{Seq})(\exists k_0)(\forall u_1 \in \text{Seq})(\exists v_1 \in \text{Seq})(\exists k_1) \cdots \\ &(\forall i) \left[E_{\langle k_0, \dots, k_{i-1} \rangle}^x \text{ is comeager in } \Sigma(\hat{s} \hat{u}_0 \hat{v}_0 \cdots \hat{u}_{i-1} \hat{v}_{i-1}) \right]. \end{aligned}$$

DEFINITION 6.2. Define an operation \mathcal{V} as follows.

$$\begin{aligned} x \in \mathcal{V}(\{E_n : n \in \omega\}) \\ \leftrightarrow (\forall u_0 \in \text{Seq})(\exists v_0 \in \text{Seq})(\exists k_0)(\forall u_1 \in \text{Seq})(\exists v_1 \in \text{Seq})(\exists k_1) \\ \dots (\forall i) \left[x \in E_{\langle \langle u_0, \langle k_0, v_0 \rangle \rangle, \dots, \langle u_{i-1}, \langle k_{i-1}, v_{i-1} \rangle \rangle} \right]. \end{aligned}$$

Clearly \mathcal{V} is positive analytical and $\mathcal{V} \sim \mathcal{A}$. Call \mathcal{V} the Vaught operation.

Define a sequence of positive analytical operations $\{S_\rho : \rho < \omega_1\}$ by the induction

$$S_0 = \mathcal{V}, \quad S_{\rho+1} = S_\rho^*.$$

If λ is limit, choose $\rho_i \uparrow \lambda$ and set for any family $\{E_n\}$,

$$\Psi(\{E_n\}) = \bigcap_{i \in \omega} \Phi_{M_{\rho_i}, M_{\rho_i}^0}(\{E_{\langle i, m \rangle} : m \in \omega\}),$$

where M_{ρ_i} is the canonical base for S_{ρ_i} . Then define $S_\lambda = \Psi^*$. It is easy to check by induction that for each ρ , $R_\rho \sim S_\rho$ and hence $\mathcal{R}^\rho = \Sigma_1^{R_\rho} = \Sigma_1^{S_\rho}$.

We shall now deduce the Category Formula of Burgess by showing that if E is in \mathcal{R}^ρ , then E^* is computed by S_ρ .

THEOREM 6.3. Let $E \subseteq \omega^\omega \times \omega^\omega$ be a set in \mathcal{R}^ρ . Then $E^{*\Sigma(s)} = \{x : E^x \text{ is comeager in } \Sigma(s)\}$ is also in \mathcal{R}^ρ .

PROOF. We shall prove the theorem by induction. Let N_ρ and M_ρ denote the canonical bases for R_ρ and S_ρ , respectively.

Assume that for all $\nu < \rho$ there are functions f_ν and g_ν such that if

$$F = R_\nu(\{F_n : n \in \omega\}),$$

with each $F_n \subseteq \omega^\omega \times \omega^\omega$ normal, then

$$F^{*\Sigma(s)} = S_\nu \left(\left\{ F_{f_\nu(n)}^{\Sigma(\hat{s}g_\nu(n))} : n \in \omega \right\} \right).$$

We shall then show that if $E = R_\rho(\{E_n\})$, then $E^{*\Sigma(s)}$ is similarly computed by S_ρ . The result then follows by observing that for each clopen E_n , E_n^* is also clopen.

Case 1. $\rho = \nu + 1$. In this case, $E = R_\nu^*(\{E_n : n \in \omega\})$. Hence by the Transfer Theorem,

$$E^x \text{ is comeager in } \Sigma(s) \leftrightarrow (\exists \eta_0 \in M_\nu)(\forall n_0 \in \eta_0)(\forall \xi_0 \in M_0)(\exists m_0 \in \xi_0) \dots$$

$$\begin{aligned} (\forall i) \left[E_{\tilde{f}_\nu(\langle \langle n_0, m_0 \rangle, \dots, \langle n_{i-1}, m_{i-1} \rangle \rangle)}^x \text{ is comeager in} \right. \\ \left. \Sigma(\hat{s}\tilde{g}_\nu(\langle \langle n_0, m_0 \rangle, \dots, \langle n_{i-1}, m_{i-1} \rangle \rangle)) \right] \end{aligned}$$

where \tilde{f}_ν and \tilde{g}_ν are related to f_ν and g_ν as in 5.4. Setting $f_\rho = \tilde{f}_\nu$ and $g_\rho = \tilde{g}_\nu$ we get

$$E^{*\Sigma(s)} = S_\rho \left(\left\{ E_{f_\rho(n)}^{\Sigma(\hat{s}g_\rho(n))} : n \in \omega \right\} \right).$$

Case 2. ρ is limit. Choose a sequence $\rho_i \uparrow \rho$ and for any normal family $\{H_n : n \in \omega\}$ set

$$H = \Phi(\{H_n\}) = \bigcap_{i \in \omega} \Phi_{N_{\rho_i}, N_{\rho_i}^0}(\{H_{\langle i, n \rangle} : n \in \omega\}).$$

Then,

$$\begin{aligned}
 H^x \text{ is comeager in } \Sigma(s) &\leftrightarrow (\forall i) \left[\left(\Phi_{N_{\rho_i} N_{\rho_i}^0}(\{ H_{\langle i, m \rangle} : m \in \omega \}) \right)^x \text{ is comeager in } \Sigma(s) \right] \\
 &\leftrightarrow (\forall i) (\exists \eta \in M_{\rho_i}) (\forall n \in \eta) (\forall \xi \in M_{\rho_i}) (\exists m \in \xi) \\
 &\quad \left[H_{\langle i, \alpha_i(\langle n, m \rangle \rangle}^x \text{ is comeager in } \Sigma(\hat{s} \beta_i(\langle n, m \rangle)) \right] \\
 &\quad \text{(by the Transfer Theorem),}
 \end{aligned}$$

where α_i and β_i are obtained as in 5.4(b). Therefore,

$$H^{\star \Sigma(s)} = \Psi \left(\{ H_{f(m)}^{\star \Sigma(s \hat{g}(m))} : m \in \omega \} \right),$$

where Ψ is the operation in 6.2 and

$$f(\langle i, n \rangle) = \langle i, \alpha_i(n) \rangle, \quad g(\langle i, n \rangle) = \beta_i(n).$$

Since $E = \Phi^*(\langle E_n \rangle)$, by applying the Transfer Theorem again we obtain the result as in Case 1.

As the R -sets have the Baire property the next result follows immediately.

COROLLARY 6.4. *If $E \in c\mathcal{R}^p$, then $E^{\star \Sigma(s)}$ is also in $c\mathcal{R}^p$.*

By repeatedly applying the Transfer Theorem and 6.3 to every level of the hierarchy of $\mathcal{B}\mathcal{R}^p$ -sets one obtains

COROLLARY 6.5. *If $E \in \mathcal{B}\mathcal{R}_\mu^p$ ($\mu, \rho < \omega_1$), then $E^{\star \Sigma(s)}$ is also in $\mathcal{B}\mathcal{R}_\mu^p$. In particular, putting $\rho = 0$, one has for any C -set E , E^* is also a C -set.*

7. The approximation theorem.

LEMMA 7.1. *For $\rho < \omega_1$, let $R_\rho = \Phi_K = R\Phi_{NN^0}$. As in the proof of 6.3 there are functions \tilde{f} and \tilde{g} such that for any normal family $\{E_n\}$ with $E = R_\rho(\langle E_n \rangle)$,*

$$\begin{aligned}
 E^x \text{ is comeager} &\leftrightarrow (\exists \eta \in \tilde{K}) (\forall n \in \eta) [E_{\tilde{f}(n)}^x \text{ is comeager in } \Sigma(\tilde{g}(n))] \\
 &\leftrightarrow (\exists \eta_0 \in M) (\forall n_0 \in \eta_0) \cdots \\
 &\quad (\forall k) \left[E_{\langle \alpha(\langle n, m \rangle) \dots \alpha(\langle n_{k-1}, m_{k-1} \rangle) \rangle}^x \text{ is comeager in} \right. \\
 &\quad \left. \Sigma(\beta(\langle n_0, m_0 \rangle) \hat{\cdot} \cdots \hat{\cdot} \beta(\langle n_{k-1}, m_{k-1} \rangle)) \right],
 \end{aligned}$$

where \tilde{f} , \tilde{g} and α, β are related as in 5.4(c) and $\Phi_M \sim \Phi_N$ and $\tilde{K} = RMM^0$.

Then, one can choose \tilde{K} such that for any $\eta \in \tilde{K}$, $\bigcup_{n \in \eta} \Sigma(\tilde{g}(n))$ is dense in ω^ω .

In fact, \tilde{K} may be taken to be the canonical base for S_ρ with the property that for any $\eta \in \tilde{K}$ and $s \in \text{Seq}$ there is an $n \in \eta$ such that $\tilde{g}(n)$ extends s .

PROOF. We shall prove this by induction on ρ . For $\rho = 0$, $\Phi_{\tilde{K}}$ is the Vaught operation \mathcal{V} . Taking \tilde{K} to be the canonical base for \mathcal{V} it is easy to check the assertion of the above theorem.

So assume $\rho > 0$ and the assertion holds for all $\nu < \rho$.

Case 1. $\rho = \nu + 1$. Then $\Phi_N = R_\nu$. Let $\eta \in \tilde{K} = RMM^0$, where M has been chosen to satisfy the assertion of the theorem for the operation R_ν .

Fix a basic clopen set $\Sigma(t)$. We have to show that there is $n \in \eta$ such that $\tilde{g}(n)$ is consistent with t (in fact, extends t).

Since $\eta \in RMM^0$,

$$(\exists \eta_0 \in M)(\forall n_0 \in \eta_0)(\forall \xi_0 \in M)(\exists m_0 \in \xi_0) \cdots \\ \cdots (\forall k)[\langle \langle n_0, m_0 \rangle \cdots \langle n_{k-1}, m_{k-1} \rangle \rangle \in \eta].$$

Fix $\eta_0 \in M$ such that for each $n_0 \in \eta_0$,

$$(\forall \xi_0 \in M)(\exists m_0 \in \xi_0)(\exists \eta_1 \in M)(\forall n_1 \in \eta_1)(\forall \xi_1 \in M)(\exists m_1 \in \xi_1) \cdots \\ (\forall k)[\langle \langle n_0, m_0 \rangle, \dots, \langle n_{k-1}, m_{k-1} \rangle \rangle \in \eta].$$

By the induction hypothesis, there is an $n_0^* \in \eta_0$ such that $g(n_0^*)$ extends t . Now pick a $\xi_0 \in M$ such that for some $m_0^* \in \xi_0$,

$$(\exists \eta_0 \in M)(\forall n_0 \in \eta_0)(\forall \xi_0 \in M)(\exists m_0 \in \xi_0) \cdots \\ \cdots (\forall k)[\langle \langle n_0, m_0 \rangle \cdots \langle n_{k-1}, m_{k-1} \rangle \rangle \in \eta].$$

Clearly, $n = \langle \langle n_0^*, m_0^* \rangle \rangle \in \eta$ and $\tilde{g}(n) = g(n_0^*) \hat{\delta}(m_0^*)$ (cf. 5.4). Thus $\tilde{g}(n)$ extends t .

Case 2. ρ is limit. Let $\rho_i \uparrow \rho$. Then $R_\rho = \Phi_N^* = R\Phi_{NN^0}$, where Φ_N is the operation given by

$$\Phi_N(\langle F_n \rangle) = \bigcap_{i \in \omega} \Phi_{N_{\rho_i} N_{\rho_i}^0}(\langle F_{\langle i, m \rangle} : m \in \omega \rangle) \quad \text{and} \quad \Phi_{N_{\rho_i}} = R_{\rho_i}.$$

In this case, if $E = \Phi_N(\langle E_n \rangle)$, then

$$E^x \text{ is comeager} \leftrightarrow (\exists \eta \in \tilde{N})(\forall n \in \eta)[E_{\tilde{f}(n)}^x \text{ is comeager in } \Sigma(\tilde{g}(n))] \\ \leftrightarrow (\forall i)(\exists \eta \in M_{\rho_i})(\forall n \in \eta)(\exists \xi \in M_{\rho_i}^0)(\forall m \in \xi) \\ [E_{\langle i, \alpha, \langle \langle n, m \rangle \rangle}^x \text{ is comeager in } \Sigma(\beta_i(\langle \langle n, m \rangle \rangle))] \\ \text{(cf. 6.3, Case 2).}$$

In view of Case 1, it suffices to show that for any $\eta \in \tilde{N}$ and for any $\Sigma(t)$, there is an $n \in \eta$ such that $\tilde{g}(n)$ extends t .

Now, \tilde{N} may be chosen such that

$$\eta \in \tilde{N} \leftrightarrow (\forall i)(\exists \eta' \in M_{\rho_i})(\forall n \in \eta')(\exists \xi' \in M_{\rho_i}^0)(\forall m \in \xi')[\langle \langle i, \langle n, m \rangle \rangle \rangle \in \eta],$$

where each M_{ρ_i} has been chosen to satisfy the assertion of the theorem.

Let $\eta \in \tilde{N}$ and fix any $i^* \in \omega$. Get $\eta^* \in M_{\rho_{i^*}}$ such that for each $n \in \eta^*$

$$(\exists \xi' \in M_{\rho_{i^*}}^0)(\forall m \in \xi')[\langle \langle i^*, \langle n, m \rangle \rangle \rangle \in \eta].$$

Since $\eta^* \in M_{\rho_{i^*}}$, by the induction hypothesis, there is an $n^* \in \eta^*$ such that $g_{i^*}(n^*)$ extends t . Get $\xi' \in M_{\rho_{i^*}}^0$ and $m^* \in \xi'$ such that $\langle \langle i^*, \langle n^*, m^* \rangle \rangle \rangle \in \eta$. Take $n = \langle \langle i^*, \langle n^*, m^* \rangle \rangle \rangle$. Then

$$\tilde{g}(n) = \beta_{i^*}(\langle \langle n^*, m^* \rangle \rangle) = g_{i^*}(n^*) \hat{\delta}_{i^*}(m^*)$$

(cf. 6.3) and hence $\tilde{g}(n)$ extends t .

LEMMA 7.2. Let \mathcal{F} be a σ -field closed under operation \mathcal{A} and let $\mathcal{B}_{\omega^\omega}$ denote the Borel σ -field on ω^ω . Let $E \subseteq \omega^\omega \times \omega^\omega$. Suppose there is a family $\{A_n\}$ such that

- (a) $A = \mathcal{A}(\{A_n\})$,
- (b) $A \subseteq E$,
- (c) each $A_n \in \mathcal{F} \otimes \mathcal{B}_{\omega^\omega}$,
- (d) A^x is comeager whenever E^x is comeager.

Then there is a set $B \in \mathcal{F} \otimes \mathcal{B}_{\omega^\omega}$ such that $B \subseteq E$ and B^x is comeager whenever E^x is comeager.

PROOF. Let $I: \omega^\omega \rightarrow \omega^\omega$ be the characteristic function of a generator for a countably generated sub- σ -field \mathcal{F}_0 of \mathcal{F} such that each $A_n \in \mathcal{F}_0 \otimes \mathcal{B}_{\omega^\omega}$. Let $I(\omega^\omega) = D$. Then, as is well known, I is a bimeasurable function between \mathcal{F}_0 and \mathcal{B}_D , the Borel σ -field on D . Set

$$\tilde{A}_n = \{(I(\alpha), \beta) : (\alpha, \beta) \in A_n\}, \quad \tilde{A} = \{(I(\alpha), \beta) : (\alpha, \beta) \in A\}.$$

Clearly, $\tilde{A}_n \in \mathcal{B}_D \otimes \mathcal{B}_{\omega^\omega}$ and $A^* = \mathcal{A}(\{\tilde{A}_n\})$. Hence \tilde{A} is an analytic set in $D \times \omega^\omega$. Get an analytic set $C \subseteq \omega^\omega \times \omega^\omega$ such that

$$\tilde{A} = C \cap (D \times \omega^\omega).$$

Then by 1.6 of [17], get $\tilde{B} \in \mathcal{A} \otimes \mathcal{B}_{\omega^\omega}$, where \mathcal{A} is the analytic σ -field on ω^ω , such that $\tilde{B} \subseteq C$ and \tilde{B}^x is comeager whenever C^x is comeager. Put

$$B = \{(\alpha, \beta) : (I(\alpha), \beta) \in \tilde{B}\}.$$

Since \mathcal{F} is closed under operation \mathcal{A} , clearly $B \in \mathcal{F} \otimes \mathcal{B}_{\omega^\omega}$ and moreover, $B \subseteq E$ and B^x is comeager whenever E^x is comeager.

We have adapted the proof of the Category Formula [2] in the next lemma.

LEMMA 7.3. Suppose we have

$$\begin{aligned} &(\forall s_0)(\exists t_0)(\forall s_1)(\exists t_1) \cdots \{(\forall a_0)(\exists b_0)(\forall a_1)(\exists b_1) \cdots P(\alpha, \beta)\} \\ &\leftrightarrow (\forall s_0)(\forall a_0)(\exists t_0)(\exists b_0)(\forall s_1)(\forall a_1)(\exists t_1)(\exists b_1) \cdots P(\alpha, \beta), \end{aligned}$$

where $\alpha = (a_0, b_0, a_1, b_1, \dots)$, $a_i, b_i \in \omega$, and $\beta = \hat{s}_0 \hat{t}_0 \hat{s}_1 \hat{t}_1 \cdots$; $s_i, t_i \in \text{Seq}$.

If \exists wins the second game, then he may do so by means of a strategy σ^* such that, to every complete play $s_0, a_0, t_0, b_0, \dots$ consistent with σ^* , there corresponds a complete play $s'_0, t'_0, s'_1, t'_1, \dots$ consistent with a winning strategy for \exists in the first game such that

$$\hat{s}_0 \hat{t}_0 \hat{s}_1 \hat{t}_1 \cdots = \hat{s}'_0 \hat{t}'_0 \hat{s}'_1 \hat{t}'_1 \cdots.$$

PROOF. Let $w_{-1}, w_0, w_1, w_2, \dots$ be an enumeration of $\omega^{<\omega}$ such that w_{-1} is the empty sequence and if w_i is an initial segment of w_j , then $i < j$. For $s \in \omega^{<\omega}$, its code, denoted by $|s|$, is its position in the enumeration.

Suppose \exists wins the second game with strategy σ . We shall now construct a winning strategy τ for player II in the first (Banach-Mazur) game. τ will be defined (by induction) in such a way that every partial play consistent with τ corresponds to a partial play consistent with σ .

Suppose $s_0, t_0, \dots, s_{n-1}, t_{n-1}$ have been defined consistent with τ and

$$\tau(s_0, t_0, \dots, s_{n-1}, t_{n-1}, s)$$

has not been defined. Let n be the code of (a_0, \dots, a_{m-1}, a) and let $\|(a_0, \dots, a_{m-1})\| = n^*$. Clearly, $n^* < n$. Let the partial play (consistent with σ) corresponding to $(s_0, t_0, \dots, s_{n^*}, t_{n^*})$ be $(a_0, u_0, \dots, a_{m-1}, u_{m-1}, b_{m-1}, v_{m-1})$ such that $\hat{s}_0 \hat{t}_0 \cdots \hat{s}_{n^*} \hat{t}_{n^*} = u_0 \hat{v}_0 \cdots \hat{u}_{m-1} \hat{v}_{m-1}$. Put

$$u = \hat{s}_{n^*+1} \hat{t}_{n^*+1} \cdots \hat{s}_{n-1} \hat{t}_{n-1} s.$$

Let

$$\sigma(a_0, u_0; b_0, v_0; \cdots; b_{m-1}, v_{m-1}, a, u) = (b, v).$$

Then

$$\tau(s_0, t_0, \dots, s_{n-1}, t_{n-1}, s) = v,$$

and the partial play associated with

$$s_0, t_0, \dots, s_{n-1}, t_{n-1}, s, v$$

is

$$a_0, u_0; b_0, v_0, \dots, b_{m-1}, v_{m-1}; a, u; b, v.$$

We shall now show that τ is a winning strategy for player II in the Banach-Mazur game. Let

$$s_0, t_0, s_1, t_1, s_2, t_2, \dots$$

be a complete play consistent with τ . We shall have to show that

$$(\forall a_0)(\exists b_0)(\forall a_1)(\exists b_1) \cdots P(\alpha, \beta),$$

where

$$\alpha = (a_0, b_0, a_1, b_1, \dots) \quad \text{and} \quad \beta = \hat{s}_0 \hat{t}_0 \hat{s}_1 \hat{t}_1 \cdots.$$

So let \forall play a_0 . Suppose $\|(a_0)\| = n_0$. By definition of τ , the partial play (consistent with σ) corresponding to $s_0, t_0, \dots, s_{n_0}, t_{n_0}$ is

$$a, \hat{s}_0 \hat{t}_0 \cdots \hat{s}_{n_0} \hat{t}_{n_0}; b_0, t_{n_0}$$

for some $b_0 \in \omega$. \exists replies with b_0 . Next suppose \forall plays a_1 and let $\|(a_0, a_1)\| = n_1$. By definition of τ , the partial play (consistent with σ) corresponding to

$$s_0, t_0, \dots, s_{n_0}, t_{n_0}, \dots, s_{n_1}, t_{n_1}$$

is

$$a_0, \underbrace{\hat{s}_0 \hat{t}_0 \cdots \hat{s}_{n_0}}_{u_0}; b_0, t_{n_0} = v_0; a_1, \underbrace{\hat{s}_{n_0+1} \hat{t}_{n_0+1} \cdots \hat{s}_{n_1}}_{u_1}; b_1, t_{n_1} = v_1.$$

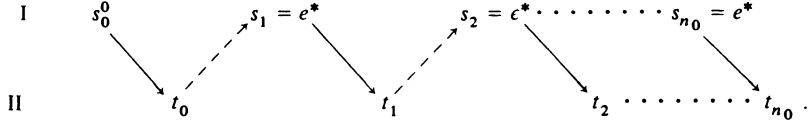
\exists replies with b_1 and the play proceeds as described. Since the play $a_0, u_0; b_0, v_0; a_1, u_1; b_1, v_1; \dots$ is consistent with σ , we have

$$P(\alpha, u_0 \hat{v}_0 u_1 \hat{v}_1 \cdots).$$

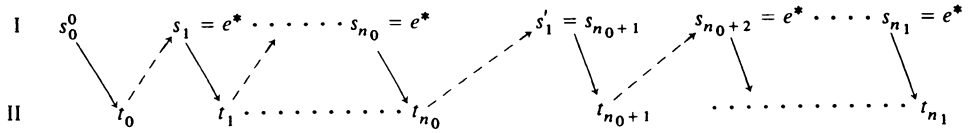
But $\beta = u_0 \hat{v}_0 u_1 \hat{v}_1 \cdots$. Hence $P(\alpha, \beta)$. Consequently, τ is a winning strategy for II in the Banach-Mazur game.

We shall now modify σ to σ^* such that any complete play consistent with σ^* corresponds to a complete play consistent with τ .

DEFINITION OF σ^* . Let \forall play a_0, s_0^0 and suppose $\|(a_0)\| = n_0$. Simulate the following partial play in the Banach-Mazur game, where II plays with strategy τ , e^* being a fixed sequence number:



According to the definition of τ , $\sigma(a_0, s_0^0 \hat{t}_0 \hat{s}_1 \hat{t}_1 \cdots \hat{s}_{n_0}) = (b_0, t_{n_0})$, for some $b_0 \in \omega$. Define $\sigma^*(a_0, s_0^0) = (b_0, t_0 \hat{s}_1 \hat{t}_1 \cdots \hat{s}_{n_0} t_{n_0})$. Next suppose \forall plays a_1, s_1' and let $\|(a_0, a_1)\| = n_1$. Consider the following partial play consistent with τ :



By definition of τ , $\sigma(a_0, s_0^0 \hat{t}_0 \cdots \hat{s}_{n_0}; b_0, t_{n_0}; a_1, \hat{s}_{n_0+1} \hat{t}_{n_0+1} \cdots \hat{s}_{n_1}) = (b_1, t_{n_1})$, $b_1 \in \omega$. Define

$$\sigma^*(a_0, s_0^0; b_0, t_0 \hat{s}_1 \hat{t}_1 \cdots \hat{s}_{n_0} t_{n_0}; a_1, s_1') = (b_1, t_{n_0+1} \hat{s}_{n_0+2} \cdots \hat{s}_{n_1} t_{n_1})$$

and so on. We shall show that σ^* is a winning strategy for \exists in the second game. First observe that any complete play

$$a_0, u_0, b_0, v_0, a_1, u_1, b_1, v_1, \dots$$

consistent with σ^* corresponds to a complete play

$$s_0, t_0, s_1, t_1, \dots$$

consistent with τ such that

$$\hat{s}_0 \hat{t}_0 \hat{s}_1 \hat{t}_1 \cdots = u_0 v_0 u_1 v_1 \cdots = \beta,$$

say. To prove that σ^* is a winning strategy, we have to show $P(\alpha, \beta)$, where $\alpha = (a_0, b_0, a_1, b_1, \dots)$. Next observe that any play consistent with σ^* is of the form

$$\begin{array}{ll} (\forall)\text{I} & a_0, s_0^0 \qquad \qquad \qquad a_1, s_1' = s_{n_0+1} \quad \cdots \\ (\exists)\text{II} & \underbrace{b_0, t_0 \hat{s}_1 \hat{t}_1 \cdots \hat{s}_{n_0} t_{n_0}} \qquad \underbrace{b_1, t_{n_0+1} \cdots \hat{s}_{n_1} t_{n_1} \cdots} \end{array}$$

such that

$$\begin{array}{ll} (\forall)\text{I} & a_0, s_0^0 \hat{t}_0 \cdots \hat{s}_{n_0} \qquad \qquad \qquad a_1, \hat{s}_{n_0+1} \hat{t}_{n_0+1} \cdots \hat{s}_{n_1} \quad \cdots \\ (\exists)\text{II} & \qquad \qquad \qquad b_0, t_{n_0} \qquad \qquad \qquad b_1, t_{n_1} \quad \cdots \end{array}$$

is consistent with σ . Consequently, we have

$$P\left(\alpha, s_0^0 \hat{t}_0 \cdots \hat{s}_{n_0} \hat{t}_{n_0} \hat{s}_{n_0+1} \hat{t}_{n_0+1} \cdots\right)$$

and hence $P(\alpha, \beta)$.

REMARK 7.3.1. The definition of σ^* is highly constructive and can be effected by a Borel function from the space of strategies to the space of strategies i.e., the map $\sigma \mapsto \sigma^*$ is Borel.

Using the above result, one immediately has the following.

COROLLARY 7.4. *If the second game in 7.3 is determined, then*

$$(\exists s_0)(\forall t_0)(\exists s_1)(\forall t_1) \cdots \{ (\exists a_0)(\forall b_0)(\exists a_1)(\forall b_1) \cdots \neg P(\alpha, \beta) \} \\ \leftrightarrow (\exists s_0)(\exists a_0)(\forall t_0)(\forall b_0)(\exists s_1)(\exists a_1)(\forall t_1)(\forall b_1) \cdots \neg P(\alpha, \beta),$$

and the conclusion of 7.3 also holds here.

REMARK. Lemma 7.3 and Corollary 7.4 yield a constructive proof of a particular case of the Game Formula of Kechris [11].

LEMMA 7.5. *Suppose $E \subseteq \omega^\omega$ be a set in \mathcal{R}^ρ , $\rho < \omega_1$ and let $R_\rho = R\Phi_N$. Assume $E = R_\rho(\langle E_n \rangle)$, where each E_n is clopen. Then,*

$$E \text{ is comeager} \leftrightarrow (\exists \eta_0 \in M)(\forall n_0 \in \eta_0)(\exists \eta_1 \in M)(\forall n_1 \in \eta_1) \cdots$$

$$(\forall k) \left[E_{\langle f(n_0), \dots, f(n_{k-1}) \rangle} \text{ is comeager in } \Sigma(g(n_0)\hat{g}(n_1)\hat{\cdots}\hat{g}(n_{k-1})) \right]$$

where f and g are some suitable functions and $\Phi_M \sim \Phi_N$. Moreover, f and g can be chosen such that, with any winning strategy σ for \exists one can associate a winning strategy σ^* such that for any run $\eta_0, n_0, \eta_1, n_1, \dots$ consistent with σ^* , the sequence $\beta = g(n_0)\hat{g}(n_1)\hat{\cdots}$ is in E .

PROOF. The first assertion follows immediately from the proof of Theorem 6.3. Moreover, by Lemma 7.1, M may be chosen such that for any $\eta \in M$ and s there is $n \in \eta$ such that $g(n)$ extends s . Therefore, we have

$$(\exists \eta_0 \in M)(\forall n_0 \in \eta_0) \cdots (\forall k) \left[E_{\langle f(n_0), \dots, f(n_{k-1}) \rangle} \text{ is comeager in } \Sigma(g(n_0)\hat{g}(n_1)\hat{\cdots}\hat{g}(n_{k-1})) \right] \\ \leftrightarrow (\exists \eta_0 \in M)(\forall n_0 \in \eta_0) \cdots (\forall k) \left[\beta \in E_{\langle f(n_0), \dots, f(n_{k-1}) \rangle} \right]$$

where $\beta = g(n_0)\hat{g}(n_1)\hat{\cdots}$. Hence,

$$(i) \quad E \text{ is comeager} \leftrightarrow (\exists \eta_0 \in M)(\forall n_0 \in \eta_0)(\exists \eta_1 \in M)(\forall n_1 \in \eta_1) \cdots \\ (\forall k) \left[g(n_0)\hat{g}(n_1)\hat{\cdots} \in E_{\langle f(n_0), \dots, f(n_{k-1}) \rangle} \right].$$

We shall prove the second assertion only for $\rho = 1$, since for higher levels the proof involves similar ideas, although notationally cumbersome.

First observe that if $F = \mathcal{A}(\langle F_n \rangle)$, with each F_n clopen, then by the Category Formula [2], we have

$$F \text{ is meager} \leftrightarrow (\forall k_0)(\forall s_0)(\exists t_0)(\forall k_1)(\forall s_1)(\exists t_1) \cdots \\ \cdots (\exists i) \left[F_{\langle k_0, \dots, k_{i-1} \rangle} \text{ is meager in } \Sigma(s_0\hat{t}_0\hat{\cdots}\hat{s}_{i-1}\hat{t}_{i-1}) \right].$$

Therefore, whenever $G = \Gamma(\langle G_n \rangle)$, where $\Gamma = \mathcal{A}^0$,

$$G \text{ is comeager} \leftrightarrow (\forall k_0)(\forall s_0)(\exists t_0)(\forall k_1)(\forall s_1)(\exists t_1) \cdots$$

$$(\exists i) \left[G_{\langle k_0, \dots, k_{i-1} \rangle} \text{ is comeager in } \Sigma(\hat{s}_0 \hat{t}_0 \cdots \hat{s}_{i-1} \hat{t}_{i-1}) \right].$$

Since $E = R\Gamma(\langle E_n \rangle)$, by 5.4.2 we have

$$(ii) \quad E \text{ is comeager} \leftrightarrow (\forall k_0^0)(\forall s_0^0)(\exists t_0^0)(\forall k_1^0)(\forall s_1^0)(\exists t_1^0) \cdots (\exists i_0) \\ (\forall k_0^1)(\forall s_0^1)(\exists t_0^1)(\forall k_1^1)(\forall s_1^1)(\exists t_1^1) \cdots (\exists i_1) \\ (\forall k_0^2)(\forall s_0^2)(\exists t_0^2)(\forall k_1^2)(\forall s_1^2)(\exists t_1^2) \cdots (\exists i_2) \cdots (\forall j) \\ \left[E_{\langle \langle k_0^0, k_1^0, \dots, k_{i_0-1}^0 \rangle, \dots, \langle k_0^{j-1}, \dots, k_{i_{j-1}-1}^{j-1} \rangle \rangle} \text{ is comeager in } \right. \\ \left. \Sigma(\hat{s}_0^0 \hat{t}_0^0 \cdots \hat{s}_{i_0-1}^0 \hat{t}_{i_0-1}^0 \cdots \hat{s}_0^{j-1} \hat{t}_0^{j-1} \cdots \hat{s}_{i_{j-1}-1}^{j-1} \hat{t}_{i_{j-1}-1}^{j-1}) \right].$$

Now, E is comeager $\leftrightarrow \Pi$ wins the game $G^{**}(E^c) \leftrightarrow (\forall s_0)(\exists t_0)(\forall s_1)(\exists t_1) \cdots [\beta \in E]$, where $\beta = \hat{s}_0 \hat{t}_0 \hat{s}_1 \hat{t}_1 \dots$. Hence (using (i)), the equivalence in (ii) reduces to

$$(iii) \quad (\forall s_0)(\exists t_0)(\forall s_1)(\exists t_1) \cdots \left\{ (\forall \alpha_0)(\exists n_0)(\forall \alpha_1)(\exists n_1) \cdots \right. \\ \left. (\forall j) \left[\beta \in E_{\langle \bar{\alpha}_0(n_0), \dots, \bar{\alpha}_{j-1}(n_{j-1}) \rangle} \right] \right\} \\ \leftrightarrow (\forall k_0^0)(\forall s_0^0)(\exists t_0^0)(\forall k_1^0)(\forall s_1^0)(\exists t_1^0) \cdots (\exists i_0) \\ (\forall k_0^1)(\forall s_0^1)(\exists t_0^1)(\forall k_1^1)(\forall s_1^1)(\exists t_1^1) \cdots (\exists i_1) \cdots (\forall j) \\ \left[\hat{s}_0^0 \hat{t}_0^0 \cdots \hat{s}_{i_0-1}^0 \hat{t}_{i_0-1}^0 \hat{s}_0^1 \hat{t}_0^1 \cdots \hat{s}_{i_1-1}^1 \hat{t}_{i_1-1}^1 \cdots \right. \\ \left. \in E_{\langle \langle k_0^0, \dots, k_{i_0-1}^0 \rangle, \dots, \langle k_0^{j-1}, \dots, k_{i_{j-1}-1}^{j-1} \rangle \rangle} \right].$$

As in [3, §11], we shall define a game G' of length ω such that \exists wins G' iff \exists wins the second game in (iii). The game G' is as follows. Though its total length is ω we think of it as consisting of *potentially* infinite sequence of subgames each consisting of *potentially* infinite sequence of rounds. If in any play of G' the j th subgame actually goes through infinitely many rounds, then the $(j+1)$ th subgame never gets started. This keeps the total length within bounds.

In the j th subgame the two players play the rounds of that subgame. The l th such round opens with \exists signalling (by a choice of 0 or 1) either a challenge or a pass. If he challenges, the whole j th subgame ends at once and the players proceed to the $(j+1)$ th subgame; in this case we record $u_j = \langle k_0^j, \dots, k_{l-1}^j \rangle$ and

$$v_j = \hat{s}_0^j \hat{t}_0^j \cdots \hat{s}_{l-1}^j \hat{t}_{l-1}^j$$

formed from the moves. If \exists passes, \forall chooses $k_l^j \in \omega$ and a sequence number s_l^j and \exists replies with $t_l^j \in \text{Seq}$; then the players proceed to the $(l+1)$ th round.

If some j th subgame goes on forever because \exists fails to challenge on any round, \exists forfeits the game. If this provision does not apply, then a sequence $(u_0, v_0; u_1, v_1, \dots)$ will have been generated. \exists wins iff for all j

$$\hat{v}_0 \hat{v}_1 \hat{v}_2 \cdots = \beta \in E_{\langle u_0, u_1, \dots, u_j \rangle}.$$

Thus the game G' is essentially of the form

$$(\forall s_0)(\forall a_0)(\exists t_0)(\exists b_0)(\forall s_1)(\forall a_1)(\exists t_1)(\exists b_1) \cdots M(\alpha, \beta),$$

where the condition M is Borel (in fact G_δ , cf. [3, §11]). We shall now make a few observations. First, observe that each subgame in the second game of (iii) is a Banach-Mazur game interlaced with operation Γ (which may be looked upon as a game!). Secondly, the condition within $[]$ is the same as the condition in the first game. Consequently, any reduction of the second game of (iii) involves the interlacing of a Banach-Mazur game with an ordinary game of length ω which is obtained by the corresponding reduction of the operation R_1 . Therefore, reducing the R_1 operation within $\{ \}$ in the first game of (iii) as above, we observe that the equivalence (iii) is equivalent to the following:

$$\begin{aligned} \text{(iv)} \quad & (\forall s_0)(\exists t_0)(\forall s_1)(\exists t_1) \cdots \{ (\forall a_0)(\exists b_0) \cdots M(\alpha, \beta) \} \\ & \leftrightarrow (\forall s_0)(\forall a_0)(\exists t_0)(\exists b_0) \cdots M(\alpha, \beta). \end{aligned}$$

Now, if \exists wins the second game in (iii), he also wins G' , the second game in (iv). By Lemma 7.3, \exists wins with a strategy σ such that for any complete play $s_0, a_0, t_0, b_0, \dots$ consistent with σ , $\beta = \hat{s}_0 \hat{t}_0 \cdots$ corresponds to a complete play of the Banach-Mazur game in (iv). Moreover, the strategy σ gives rise to a strategy σ^* in the second game of (iii) such that, if $s_0^0, t_0^0, k_0^0, \dots, i_0; s_0^1, t_0^1, k_0^1, \dots, i_1; \dots$ is a complete run consistent with σ^* , then the sequence $\beta = s_0^{0^*} t_0^{0^*} \cdots s_{i_0-1}^{0^*} t_{i_0-1}^{0^*} s_0^{1^*} t_0^{1^*} \cdots$ is also produced by a complete run of the game G' , when \exists plays with σ . Consequently, β satisfies the condition within $\{ \}$ in (iv) and as observed above, $\beta \in E$.

REMARK. For sets E at the higher level we have an equivalence similar to (iii). On the left side of the equivalence we have a Banach-Mazur game followed by the corresponding R -operation, and on the right side we have Banach-Mazur games "interlaced" with the operation (regarded as a game played on ω) with dummy moves, if necessary. As in the proof above, the second game can be reduced to a game of length ω (see [3]) and the proof proceeds exactly as above.

REMARKS 7.6 (ON THE Δ -TRANSFORM). We shall now look at the dual of the $*$ -transform and obtain some analogous results. The proofs being very similar to what we have already done, we shall omit the details.

For a set $E \subseteq X \times Y$ and $U \subseteq Y$, put $E^{\Delta U} = \{ x \in X : E^x \text{ is nonmeager in } U \}$.

Note that Δ is the dual of the $*$ -transform in the sense that $E^{\Delta U} = [(E^c)^* U]^c$. The Δ -transform behaves very much like the $*$ -transform and we have the counterpart of Lemma 5.3:

With notation as in Lemma 5.3, we have

$$E^{\Delta \Sigma(s)} = \bigcap_{\mu < \omega_1} [E_e^\mu]^{\Delta \Sigma(s)} = \bigcup_{\mu < \omega_1} [E_e^\mu - T^\mu]^{\Delta \Sigma(s)}.$$

This decomposition of the set E^Δ suggests, as in the case of E^* , that $(E^\Delta)^c$ may be obtained as a fixed point of an inductive operator. Indeed we have the following counterpart of Theorem 5.4.

THEOREM. Let Φ_N and Φ_M be two $\delta - s$ operations such that Φ_N preserves the Baire property and Φ_M is normal and subsumes both $(\text{countable}) \cup$ and \cap . Suppose there are functions f_1 and g_1 such that for any normal family $\{E_p\}$ of subsets of $\omega^\omega \times \omega^\omega$ with $E = \Phi_N(\{E_p\})$,

$$E^{\Delta\Sigma(s)} = \Phi_M\left(\left\{E_{f_1(p)}^{\Delta\Sigma(\hat{s}g_1(p))} : p \in \omega\right\}\right).$$

Then for any regular normal family $\{F_p\}$, if $F = R\Phi_N(\{F_p\})$ then

F^x is nonmeager in $\Sigma(s)$

$$\leftrightarrow (\exists \eta_0 \in M')(\forall n_0 \in \eta_0)(\exists \eta_1 \in M')(\forall n_1 \in \eta_1) \cdots$$

$$(\forall k) \left[F_{\langle f'(n_0), \dots, f'(n_{k-1}) \rangle}^x \text{ is nonmeager in } \Sigma(\hat{s}g'(n_0) \cdots \hat{s}g'(n_{k-1})) \right],$$

where f' and g' are suitable functions (independent of $\{F_p\}$) and $\Phi_{M'} = \Phi_M \cup \cap$.

PROOF. First observe that for any regular family $\{F_p\}$, $F_s^\mu \subseteq F_t$ if $t < s$. As in Theorem 5.4(c) we define a set relation operative on ω as follows:

$$t \in \Gamma_x(A) \leftrightarrow F_{(t)_0}^x \text{ is meager in } \Sigma((t)_1) \vee$$

$$(\forall \eta \in M)(\exists n \in \eta)(\forall u)(\exists v) \left[\langle (t)_0 \langle f_1(n) \rangle, (t)_1 g_1(n) \hat{u} v \rangle \in A \right].$$

To obtain the result, one then shows that $x \in F^{\Delta\Sigma(s)} \leftrightarrow \langle e, s \rangle \notin \Gamma_x^\infty$. As in 6.3, this theorem immediately implies that for a set $E \in \mathcal{R}^p$, if $E = R\Phi_N(\{E_n\})$, then

$$(9) \quad x \in E^\Delta \leftrightarrow (\exists \eta_0 \in M)(\forall n_0 \in \eta_0)(\exists \eta_1 \in M)(\forall n_1 \in \eta_1) \cdots$$

$$(\forall k) \left[E_{\langle f'(n_0), \dots, f'(n_{k-1}) \rangle}^x \text{ is nonmeager in } \Sigma(g'(n_0) \cdots g'(n_{k-1})) \right],$$

for suitable M, f', g' . Moreover, one can choose M such that

$$\eta \in M^0 \& s \in \text{Seq} \rightarrow (\exists n \in \eta)(g'(n) \text{ extends } s).$$

Finally, observe that by (9) we have for $E \subseteq \omega^\omega$

$$(v) \quad E \text{ is meager} \leftrightarrow (\exists \eta_0 \in M^0)(\forall n_0 \in \eta_0)(\exists \eta_1 \in M^0)(\forall n_1 \in \eta_1) \cdots$$

$$(\exists k) \left[g'(n_0) \hat{g}'(n_1) \cdots \notin E_{\langle f'(n_0), \dots, f'(n_{k-1}) \rangle} \right].$$

Arguing as in Lemma 7.5 and invoking Corollary 7.4, we may show that \exists can win the game (v) (if he does so) by a strategy σ^{**} such that if $\eta_0, n_0, \eta_1, n_1, \dots$ is a complete play consistent with σ^{**} , then $\beta = g'(n_0) \hat{g}'(n_1) \cdots$ is not in E .

The next theorem is the first step towards our approximation theorem.

THEOREM 7.7. Let $E \subseteq \omega^\omega \times \omega^\omega$ be a set in \mathcal{R}^p ($p > 0$). Then there is a set $B \in \mathcal{B}\mathcal{R}_0^p \otimes \mathcal{B}_{\omega^\omega}$ such that $B \subseteq E$ and B^x is comeager on E^* .

PROOF. Let $R_p = R\Phi_N$ and assume $E = R\Phi_N(\{E_n\})$, with each E_n clopen. As in 6.3, there are functions f and g such that

$$(i) \quad x \in E^* \leftrightarrow (\exists \eta_0 \in M)(\forall n_0 \in \eta_0)(\exists \eta_1 \in M)(\forall n_1 \in \eta_1) \cdots$$

$$(\forall k) \left[E_{\langle f(n_0), \dots, f(n_{k-1}) \rangle}^x \text{ is comeager in } \Sigma(g(n_0) \cdots g(n_{k-1})) \right]$$

where M is such that $\Phi_M \sim \Phi_N$. Further, by Theorem 5.1, there is a \mathcal{BR}_0^p -measurable function $x \mapsto \tau_x$ such that whenever $x \in E^*$, τ_x is a winning strategy for the player \exists in the game (i). Moreover, by Lemma 7.5 and Remark 7.3.1, we may assume that the function $x \mapsto \tau_x$ is such that for any complete play $\eta_0, n_0, \eta_1, n_1, \dots$ agreeing with τ_x the sequence $\beta = g(n_0)\hat{g}(n_1)\hat{g}(n_2)\hat{\cdot} \dots$ is in E^* , and that the base M has been chosen as in Lemma 7.1. Define for each $k \geq 1$,

$$W_k(s, x) \leftrightarrow \text{Seq}(s) \& \text{lh}(s) = k \& (\forall i < \text{lh}(s))[(s)_i \in \tau_x(s \upharpoonright i)];$$

i.e., W_k^x consists of the codes of the first k possible moves of \forall when the existential player plays according to the strategy τ_x . Clearly, $W_k \in \mathcal{BR}_0^p$. Define

$$C(x, y) \leftrightarrow (\exists \alpha)(\forall k)[W_k(\bar{\alpha}(k), x) \& y \in \Sigma(g(\alpha(0))\hat{\cdot} \dots \hat{\cdot} g(\alpha(k-1)))] \& x \in E^*.$$

Plainly, C is the result of operation \mathcal{A} on sets in $\mathcal{BR}_0^p \otimes \mathcal{B}_{\omega^\omega}$. We shall show that (a) $C \subseteq E$ and (b) C^x is comeager on E^* .

If $(x, y) \in C$ then $x \in E^*$ and \exists wins the game (i). Further, there is a sequence $\{n_k : k \in \omega\}$ such that n_0, n_1, n_2, \dots are the moves of \forall when \exists plays according to τ_x and $y = g(n_0)\hat{g}(n_1)\hat{\cdot} \dots$. Consequently, $y \in E^*$. To prove (b), fix an x such that E^x is comeager. Then

$$C^x = \{y | (\exists \alpha)(\forall k)[W_k(\bar{\alpha}(k), x) \& y \in \Sigma(g(\alpha(0))\hat{\cdot} \dots \hat{\cdot} g(\alpha(k-1)))]\}.$$

Define

$$A_{\langle n_0, \dots, n_{k-1} \rangle} = \begin{cases} \Sigma(g(n_0)\hat{g}(n_1)\hat{\cdot} \dots \hat{\cdot} g(n_{k-1})) & \text{if } W_k(\langle n_0, \dots, n_{k-1} \rangle, x), \\ \emptyset & \text{otherwise.} \end{cases}$$

Then, $C^x = \mathcal{A}(\langle A_i \rangle)$. Hence to show C^x is comeager, it is enough to show (by Vaught's Formula) that

$$(\forall s_0)(\exists k_0)(\exists t_0)(\forall s_1)(\exists k_1)(\exists t_1) \dots \\ (\forall i) \left[A_{\langle k_0, \dots, k_{i-1} \rangle} \text{ is comeager in } \Sigma(\hat{s}_0 \hat{t}_0 \hat{\cdot} \dots \hat{\cdot} \hat{s}_{i-1} \hat{t}_{i-1}) \right].$$

So let \forall play s_0 . By Lemma 7.1, pick $k_0 \in \tau_x(\langle \rangle)$ such that $g(k_0)$ extends s_0 . Then \exists replies with k_0 and a sequence number t_0 such that $\hat{s}_0 \hat{t}_0 = g(k_0)$. Next when \forall plays s_1 , \exists plays $k_1 \in \tau_x(\langle k_0 \rangle)$ such that $g(k_1)$ extends s_1 and then plays t_1 such that $\hat{s}_1 \hat{t}_1 = g(k_1)$, and so on. By adopting this strategy, \exists ensures that for each i ,

$$A_{\langle k_0, \dots, k_{i-1} \rangle} \text{ is comeager in } \Sigma(\hat{s}_0 \hat{t}_0 \hat{\cdot} \dots \hat{\cdot} \hat{s}_{i-1} \hat{t}_{i-1}).$$

Thus, C^x is comeager. An application of Lemma 7.2 gives the required set B .

The next lemma is the counterpart of Theorem 7.7.

LEMMA 7.8. *Let $E \subseteq \omega^\omega \times \omega^\omega$ be a set in \mathcal{R}^p . Let $E^\# = \{x : E^x \text{ is meager}\}$. Then $E^\# \in c\mathcal{R}^p$ and there is a set $C \in \mathcal{BR}_0^p \otimes \mathcal{B}_{\omega^\omega}$ (in fact, in $\sigma(\mathcal{R}^p) \otimes \mathcal{B}_{\omega^\omega}$) such that $E \subseteq C$ and C^x is meager on $E^\#$.*

PROOF. Plainly, $E^\# \in c\mathcal{R}^p$, by Corollary 6.4. Let $f, g, N, M, \{E_n\}$ be as in Remark 7.6 with $\{E_n\}$ regular. Then, by (v) of Remark 7.6,

$$(i) \quad E^x \text{ is meager} \leftrightarrow (\exists \eta_0 \in M^0)(\forall n_0 \in \eta_0)(\exists \eta_1 \in M^0)(\forall n_1 \in \eta_1) \dots \\ (\exists k) \left[g(n_0)\hat{g}(n_1)\hat{\cdot} \dots \notin E_{\langle f(n_0), \dots, f(n_{k-1}) \rangle}^x \right].$$

Further, by Remark 5.1.1 there is a $\sigma(\mathcal{R}^p)$ -measurable function $x \mapsto \tau_x$ such that τ_x is a winning strategy for \exists in the above game whenever $x \in E^\#$. Moreover, τ_x can be chosen as in Remark 7.6. Call $u = \langle n_0, \dots, n_{k-1} \rangle$ good with respect to τ_x and β if β extends $g(n_0) \hat{\cdot} \dots \hat{\cdot} g(n_{k-1})$ and n_0, \dots, n_{k-1} are the first k possible moves of \forall when \exists plays with strategy τ_x . Define

$$\begin{aligned} T(u, x, \beta) &\leftrightarrow \text{Seq}(u) \& (\forall i < \text{lh}(u)) [(u)_i \in \tau_x(u \upharpoonright i)] \\ &\& \beta \in \Sigma(g((u)_0) \hat{\cdot} g((u)_1) \hat{\cdot} \dots \hat{\cdot} g((u)_{\text{lh}(u)-1})) \\ &\& (\forall n \in \tau_x(u)) [\beta \notin \Sigma(g((u)_0) \hat{\cdot} g((u)_1) \hat{\cdot} g((u)_2) \hat{\cdot} \dots \hat{\cdot} g((u)_{\text{lh}(u)-1}) \hat{\cdot} g(n))]. \end{aligned}$$

In other words, for each x, β , the section $T_{x,\beta}$ consists of all maximal good sequences with respect to τ_x and β . Clearly, $T \in \sigma(\mathcal{R}^p) \otimes \mathcal{B}_{\omega^\omega}$ and, moreover, $T^{u,x}$ is closed nowhere-dense for each good sequence u and x . Now define

$$C'(x, \beta) \leftrightarrow x \in E^\# \& (\exists u) T(u, x, \beta).$$

Set

$$C = C' \cup (\omega^\omega | E^\#) \times \omega^\omega.$$

Since $E^\# \in c\mathcal{R}^p$, $C \in \sigma(\mathcal{R}^p) \otimes \mathcal{B}_{\omega^\omega}$; and to conclude the proof we shall show that $E \subseteq C$. So let $(x, \beta) \in E$. If $x \notin E^\#$, then we are done. So assume $x \in E^\#$. Therefore, \exists wins the game (i) with strategy τ_x . Now, suppose when \exists plays the game (i) according to τ_x , \forall is able to play n_0, n_1, n_2, \dots such that for each k , $\langle n_0, \dots, n_{k-1} \rangle$ is a good sequence with respect to τ_x and β . Then

$$\beta = g(n_0) \hat{\cdot} g(n_1) \hat{\cdot} \dots$$

and hence, by Remark 7.6, $\beta \notin E^x$. Consequently, since we have $\beta \in E^x$, \forall is able to play n_0, n_1, \dots 'consistent' with β only up to a finite stage. Let u be the code of the maximal sequence. Then $T(u, x, \beta)$ and hence $(x, \beta) \in C$.

The following gives the approximation at the first level.

LEMMA 7.9. *Let $E \subseteq \omega^\omega \times \omega^\omega$ be a set in \mathcal{R}^p ($p > 0$). Then there are sets B, C in $\mathcal{B}\mathcal{R}_0^p \otimes \mathcal{B}_{\omega^\omega}$ such that $B \subseteq E \subseteq C$ and $C^x - B^x$ is meager for each x .*

PROOF. Let $T_s = E^{\star \Sigma(s)}$; $s \in \text{Seq}$. As E^x satisfies the Baire property for each x ,

$$\bigcup T_s = \{x : E^x \text{ is nonmeager}\}.$$

Since each $\Sigma(s)$ is a (recursive) homeomorphic copy of ω^ω , we can apply Theorem 7.7 to get B_s as in the theorem. Put

$$B = \bigcup B_s \cap (T_s \times \omega^\omega).$$

Clearly, $B \subseteq E$, $B \in \mathcal{B}\mathcal{R}_0^p \otimes \mathcal{B}_{\omega^\omega}$ and for each x , $E^x - B^x$ is meager.

To get the set C , work with E^c and argue as above using Lemma 7.8.

REMARK. Lemmas 7.7–7.9 are the analogues of Lemmas 1.3, 1.5 and 1.6 of Srivatsa for C -sets [17]. Naturally, from now on it is going to be a repetition of Srivatsa's techniques.

APPROXIMATION THEOREM 7.10. *Let A be an \mathcal{R}_μ^ρ subset of $\omega^\omega \times \omega^\omega$ ($\rho > 0$, $\rho, \mu < \omega_1$). Then there are sets B and C in $\mathcal{B}\mathcal{R}_\mu^\rho \otimes \mathcal{B}_{\omega^\omega}$ such that $B \subseteq A \subseteq C$ and $C^x - B^x$ is meager for each x .*

Moreover, if $A \in \mathcal{B}\mathcal{R}_\mu^\rho$, then one can find B, C in $\mathcal{B}\mathcal{R}_\mu^\rho \otimes \mathcal{B}_{\omega^\omega}$ with the above properties.

PROOF. We shall prove this by induction on μ . When $\mu = 0$, this is just Lemma 7.9. So suppose $\mu > 0$ and that the result is true for all $\nu < \mu$. Let $A \in \mathcal{R}_\mu^\rho$ and let $R_\rho = R\Phi_N$. Then A is the result of operation R_ρ on a family $\{A_n\}$ of sets from $c[\bigcup_{\nu < \mu} \mathcal{R}_\nu^\rho]$. By the induction hypothesis, for each n we have B_n and C_n both in $\bigcup_{\nu < \mu} \mathcal{B}\mathcal{R}_\nu^\rho \otimes \mathcal{B}_{\omega^\omega}$ such that $B_n \subseteq A_n \subseteq C_n$ and $C_n^x - B_n^x$ is meager for each x . Let

$$\tilde{B} = R_\rho(\{B_n\}) \quad \text{and} \quad \tilde{C} = R_\rho(\{C_n\}).$$

Then, $\tilde{B} \subseteq A \subseteq \tilde{C}$ and for each x , $\tilde{C}^x - \tilde{B}^x \subseteq \bigcup_n (C_n^x - B_n^x)$, and hence meager. To complete the proof it is enough to get $B \subseteq \tilde{B}$ and $C \supseteq \tilde{C}$ such that B and C are in $\mathcal{B}\mathcal{R}_\mu^\rho \otimes \mathcal{B}_{\omega^\omega}$, $\tilde{B}^x - B^x$ is meager and $C^x - \tilde{C}^x$ is meager for each x . We will show how to obtain B ; C can be obtained similarly.

Let \mathcal{F} be the σ -field generated by $\bigcup_{\nu < \mu} \mathcal{B}\mathcal{R}_\nu^\rho$. Then $\tilde{B} = R_\rho(\{B_n\})$ with each B_n in $\mathcal{F} \otimes \mathcal{B}_{\omega^\omega}$. Thus, one can obtain a countably generated sub- σ -field \mathcal{G} of \mathcal{F} such that each $B_n \in \mathcal{G} \otimes \mathcal{B}_{\omega^\omega}$. Fix a countable generator of \mathcal{G} and let $f: (\omega^\omega, \mathcal{G}) \rightarrow \omega^\omega$ be its characteristic function. Put $M = f(\omega^\omega)$. Then $(\omega^\omega, \mathcal{G})$ and (M, \mathcal{B}_M) are Borel isomorphic. For each n , let $B'_n = \{(f(x), y) : (x, y) \in B_n\}$. Then $B'_n \in \mathcal{B}_{M \times \omega^\omega}$, for each n . Let $D = \{(f(x), y) : (x, y) \in \tilde{B}\}$. Then $D = R_\rho(\{B'_n\})$. Hence, there is a set $E \subseteq \omega^\omega \times \omega^\omega$ in \mathcal{R}^ρ such that $E \cap (M \times \omega^\omega) = D$. Apply 7.9 and get $\tilde{E} \subseteq E$ such that $\tilde{E} \in \mathcal{B}\mathcal{R}_0^\rho \otimes \mathcal{B}_{\omega^\omega}$ and $E^x - (\tilde{E})^x$ is meager for each x . Let $\tilde{f}: \omega^\omega \times \omega^\omega \rightarrow \omega^\omega \times \omega^\omega$ be the map

$$\tilde{f}(x, y) = (f(x), y).$$

Put $B = (\tilde{f})^{-1}(\tilde{E})$. Note that since f is a bimeasurable map of $(\omega^\omega, \mathcal{G})$ and (M, \mathcal{B}_M) ,

$$f^{-1}(\mathcal{R}^\rho \upharpoonright \omega^\omega) \subseteq \{A \subseteq \omega^\omega : A \text{ is the result of operation } R_\rho \text{ on sets in } \mathcal{G}\} \subseteq \mathcal{R}_\mu^\rho.$$

Consequently, $f^{-1}(\mathcal{B}\mathcal{R}_0^\rho \upharpoonright \omega^\omega) \subseteq \mathcal{B}\mathcal{R}_\mu^\rho$. Thus $B \in \mathcal{B}\mathcal{R}_\mu^\rho \otimes \mathcal{B}_{\omega^\omega}$ and clearly, $\tilde{B}^x - B^x$ is meager for each x .

The second assertion follows from the first by observing that the class of all sets for which the result holds is closed under Φ_N and complementation.

The next proposition is quite well known and follows from the Von Neumann selection theorem.

PROPOSITION. *Let (T, \mathcal{M}) be a measurable space, \mathcal{M} being a σ -field closed under operation \mathcal{A} , and let Y be a Polish space. Let $B \in \mathcal{M} \otimes \mathcal{B}_Y$ have nonempty vertical sections. Then B has an \mathcal{M} -measurable selection.*

As a consequence of the approximation theorem we have the following

THEOREM 7.11. *Suppose $A \subseteq \omega^\omega \times \omega^\omega$ is a $\mathcal{B}\mathcal{R}_\mu^\rho$ set ($\rho > 0$) such that A^x is nonmeager for each x . Then A has a $\mathcal{B}\mathcal{R}_\mu^\rho$ -measurable selection.*

PROOF. By Theorem 7.10, get $B \in \mathcal{BR}_\mu^p \otimes \mathcal{B}_{\omega^\omega}$ such that $B \subseteq A$ and $A^x - B^x$ is meager for each x . B^x is then nonempty for each x , and the above proposition yields the result.

The next selection theorem is due to Burgess [3].

THEOREM 7.12. *Let $F: \omega^\omega \rightarrow \omega^\omega$ be a multifunction such that $F(x)$ is nonmeager in its closure $\text{cl}(F(x))$. If F is \mathcal{BR}^p -measurable and its graph $\text{Gr}(F) \in \mathcal{BR}^p$, then F has a \mathcal{BR}^p -measurable selection.*

PROOF. Define G by $G(x) = \text{cl}(F(x))$. Then G is a closed-valued, \mathcal{BR}^p -measurable multifunction. Hence there is a map $g: \omega^\omega \times \omega^\omega \rightarrow \omega^\omega$ such that g is $\mathcal{BR}^p \otimes \mathcal{B}_{\omega^\omega}$ -measurable and $g(x, \cdot)$ is continuous, open and onto $G(x)$, for each x (cf. [18]). Define $G' \subseteq \omega^\omega \times \omega^\omega$ by

$$G' = \{ (x, y) : g(x, y) \in F(x) \}.$$

As $\text{Gr}(F) \in \mathcal{BR}^p$ and g is $\mathcal{BR}^p \otimes \mathcal{B}_{\omega^\omega}$ -measurable, G' is in \mathcal{BR}^p . Also, as the inverse image of a nonmeager set under a continuous open map is nonmeager, G' has nonmeager sections. By Theorem 7.11, G' has a \mathcal{BR}^p -measurable selection g' . Then $f(x) = g(x, g'(x))$ is a \mathcal{BR}^p -measurable selection for F .

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